VOT NUMBER: 79107

Development of Appointment Scheduling Agent Using Distributed Constraint Satisfaction (DisCS) PEMBANGUNAN AGEN PENJADUALAN TEMUJANJI MENGGUNAKAN KEPUASAN KEKANGAN TERAGIH (DisCS)

NURULHUDA FIRDAUS MOHD AZMI ZETI DARLEENA ERI NUZULHA KHILWANI IBRAHIM ROZANA DIANA AHMAD RUSLI

RESEARCH VOTES NUMBER: 79107

Pusat Pengurusan Penyelidikan Universiti Teknologi Malaysia

2008

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DEDICATION

In the name of ALLAH the most merciful and compassionate. Salawat and salaam to our great prophet Muhammad SAW and his family and friends. May ALLAH bestow His blessing upon their souls. All the praises and gratefulness to ALLAH of because for His blessing and willing may this paper be presented today.

Our highest appreciation dedicated to the Ministry of Scien ce, Technology and Innovation (MOSTI) for the grant that was made available for us to carry out this study.

Our big thank you to Centre of Research Management, UTM for the support and guidance in managing this research grant.

To all the researchers, Pn Zeti Darleena bt Eri, Cik Nuzulha Khilwani Bt Ibrahim and Cik Rozana Diana Bt Ahmad Rusli, thank you very much for all the fruitfull discussions, your hard work, perseverance and the support in completing this study.

Last but not least, our gratitude goes to all the staffs of Centre for Advanced Software Engineering (CASE) at UTM City Campus for providing us with all the facilities in ensuring the smoothness of our work.

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ABSTRACT

Development of Appointment Scheduling Agent Using Distributed Const raint

Satisfaction (DisCS)

(Keywords: Distributed Artificial Intelligence, Distributed Computing System (DCS), Distributed Constraint Satisfaction (DisCS), Service Operation Scheduling, Multiple

Agent)

Scheduling problem (that is assigning resources and t ime points to given tasks) arise in many real-world domains, for example reservation environment. A reservation environment typically refers to a parallel -machine with n jobs available for processing. The processing time of each job has to fit within a giv en time window and with n jobs available for processing. The goal of this research is to develop appointment scheduling agent in reservation environment with implementation of Distributed Constraint Satisfaction (DisCS). The aim of this development is to optimize the reservation schedule towards efficiency schedule. In the other hand, to improve a better operation system level in organization. In this study, several algorithms were produced for different types of appointment policies and constraints. This study follows the Rational Unified Process (RUP) methodology as software development life cycle (SDLC) in the software development process. The design of the system architecture is illustrated using Unified Modeling Language (UML) technique.

Key Researchers:

Nurulhuda Firdaus Mohd Azmi (Project Leader) Zeti Darleena Eri Nuzulha Khilwani Ibrahim Rozana Diana Ahmad Rusli

> Email: <u>huda@utm.my</u> Phone Num.: 03-26154901 Vote Num: 79107

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ABSTRAK

Pembangunan Agen Penjadualan Temujanji Menggunakan Kepuasan Kekangan Teragih (DisCS)

(Keywords: Distributed Artificial Intelligence, Distributed Computing System (DCS), Distributed Constraint Satisfaction (DisCS), Service Operation Scheduling, Multiple

Agent)

Masalah penjadualan didefinasikan sebagai kesulitan untuk menguruskan sumber dan masa untuk satu-satu tugasan. Ianya wujud pada berbagai situasi di tempat kerja, sebagai contoh dalam tugas penempahan. Suasana tugas penempahan biasanya merujuk kepada suatu alat pemprosesan yang selari dengan bilangan (n) tugasan yang perlu diproses. Masa pemprosesan untuk satu-satu tugasan haruslah menepati satu masa yang ditentukan supaya kesemua bilangan (n) tugasan yang ada dapat diproses dalam satu jumlah masa keseluruhan yang diperuntukan. Matlamat kajian ini adalah untuk membangunkan agen penjadualan temujanji untuk diguna pakai dalam suasana tugas penempahan dengan menggunakan kaedah Kepuasan Kekangan Teragih. Kajian ini bertujuan untuk mengoptimakan jadual penempahan kepad a satu jadual yang lebih efisien, atau dalam erti kata lain, untuk memperbaiki dan meningkatkan tahap sistem operasi dalam organisasi. Dalam kajian ini beberapa algoritma telah dihasilkan untuk berbagai jenis polisi penempahan dan kekangannya. Kajian ini a dalah menuruti metodologi RUP sebagai kitar hayat pembangunan perisian dalam proses pembangunan perisian. Rekabentuk sistem dibangunkan menggunakan teknik Unified Modeling Language (UML).

Penyelidik Utama:

Nurulhuda Firdaus Mohd Azmi (Ketua Projek) Zeti Darleena Eri Nuzulha Khilwani Ibrahim Rozana Diana Ahmad Rusli

> Emel: <u>huda@utm.my</u> No. Tel.: 03-26154901 No. Vot: 79107

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ABBREVIATIONS

ASRS	Appointment Scheduling for Reservation System
RUP	Rational Unified Process
SDLC	Software Development Life Cycle
UML	Unified Modeling Language
VIC	Vehicle Inspection Center

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A Gantt Chart

CHAPTER 1

INTRODUCTION

1.1 Introduction

Scheduling problem (that is assigning resources and time points to given tasks) arise in many real-world domains, for example reservation environment. A reservation environment typically refers to a parallel -machine with n jobs available for processing. The processing time of each job has to fit within a given time window and with n jobs available for processing. The processing time of each job has to fit within a given time window and there may or may not be a ny slack. The time window of job is specified by a release date rj and a due date dj.

It may be the case that is not possible to process all n jobs, and the scheduler has to decide which jobs to process. Several objectives may be considered, for example, maximizing the number of job processed or maximizing the total amount of processing. Recently, several researchers have applied agent technology to such problems especially in manufacturing and business process control. Because of the widely spread lack of theoretically foundation for this kind of collaborative problem solving, the constraint satisfaction/optimization community has recently extended the theoretically well -founded constraint satisfaction problem (CSP) model to a distributed constraint satis faction problems (DCSP).

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A distributed CSP is a constraint satisfaction problem in with variables and constraints are distributed among multiple agents. The variables are connected by constraints that define the constraints network among the agents.

As a result, the search algorithm for solving these problems is distributed algorithm. The algorithm is run by agents that communicate by sending and receiving messages. In general, messages contain information about assignments of values to variables and refutation of assignments, by agents that have no compatible assignments, to their own variables.

1.2 Research Objectives

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The goal of this research is to develop appointment scheduling agent in reservation environment with implementation of Distributed Constraint Satisfaction (DisCS). In order to arrive this goal there are other objectives that need to be achieved, which are:

- To develop applicable mathematical model based on Distributed Constraint Satisfaction (DisCS) in appointment schedule.
- 2. To design and develop system architecture of appointment scheduling agent.

1.3 Background of Case Study – Vehicle Inspection Centre

The number of vehicles undergoing inspection appears to be gradually increasing every year. The current workload based on 2004 statistics indica tes that the VIC performs 3.2 million commercial tests per year with 750,000 vehicles on the road. Within this statistic, VIC outstations in the rural areas service some 5%, a mobile station which covers three locations, services about 2% (i.e. around 900 – 1000 vehicles per month for each location) and the rest are carried by VIC fixed centers. All the VIC's branches have their similar computerized inspection system. There are statistics of inspections for three highest inspection volumes by August 2004.

- i. Wangsa Maju : 17, 091 vehicles inspected
- ii. Padang Jawa : 15, 613 vehicles inspected
- iii. Johor Bahru : 15, 306 vehicles inspected

The other VIC branches are Cheras, B.Maluri, Banting, Ipoh, Taiping, Teluk Intan, Mak Mandin, Teluk Kumbar, Sungai Petani, Alor Setar, Arau, Kuantan, Kota Bharu, Kuala Terengganu, Raub, Muar, Kluang, Alor Gajah, Seremban, Kota Kinabalu, Tawau, Sandakan, Kuching, Sibu, Miri, Labuan, Pelabuhan Klang, Kuala Krai and Kijal. Total inspections from all the centers for August 2004 is 177, 950 vehicles.

Generally, the inspection services offered at the VIC are as follows:

i.	Initial Inspection	: Inspection to determine vehicle status before
		registration with the Road Transport Department, or
		before transfer of ownership for commercial
		vehicles.
ii.	Routine Inspection	: Routine half-yearly checks to gauge
		roadworthiness of commercial vehicles and ensure
		compliance with Malaysia Construction and Use
		Rules 1959.
iii.	Re-Inspection	: To be conducted after a vehicle fails initial or
		routine inspection.
iv.	Special Inspection	: To determine roadworthiness of modified vehicles,
		as well as verification of imported vehicles.
v.	Accident Inspection	: In aid of police investigations of fatal accidents
		involving vehicles.

There are 10 points of vehicle inspection in the centre and it is defined as follow:

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- i. Identification check verifies the engine and chassis number of the vehicle to be tested.
- Above carriage check the vehicle is then scrutinized for surface defects and non-compliance with road transport act 1987.

- Emission test density of particles in diesel engines and volume of gaseseg HC, CO etc in petrol engines are gauged through free acceleration tests.
- iv. Side-slip test lateral movements of the vehicle are tested for optimum road-handling criteria.
- v. Suspension test checks the effectiveness of the suspension system (springs, shock absorber, and joints) on each axle.
- vi. Brake test the various brake efficiencies, such as service brake efficiency, dynamic unbalance, run-out, residual force and parking brake are evaluated for maximum performance.
- vii. Speedometer checks the road speed at the wheels against the value indicated on the vehicle's speedometer.
- viii. Headlight test measures the intensity and projection of the vehicle's headlights
- ix. Undercarriage check a vehicle's undercarriage is thoroughly scrutinized for defects while it is tested on an axle play detector.
- Computerized analysis a state-of-the-art central computer controls and monitors each test phase. All relevant test data is then analyzed and the final result printed out.

1.3.1 Johor Bahru Vehicle Inspection Centre

A VIC in the southern part of peninsular Malaysia, named VIC Johor Bahru (JB) has been selected as a case study for this research. It was identified by the VIC top management as one of the busiest centres and is known to have the highest level of operational problems. VIC JB is situated in an industrial area with a high volume of vehicles coming for inspection. On average, the number of inspected vehicles in VIC JB is 400 per day and is increase by an average of 5% every year. VIC JB is led by a Centre Manager who is assisted by two supervisors and a quality control officer. There are currently a total of 23 vehicle examiners at this site. The size of the VIC premises is about 7.2 km² (i.e. 120 metres X 60 metres) and is an open area. The layout of VIC JB is as in **Figure 1.1**. Queuing areas for light and heavy vehicles are located separately, except for initial inspections. The separation is to ease movement and for safety purposes. As an addition, separa te queuing areas for light and heavy vehicles are the reason of why the VIC has two entrances and one exit. Two entrances, each for light and heavy vehicle are designed at two different locations as shown in **Figure 1.1**, each of which located near its queue lane. This is to enable smooth flow for both types of vehicles.

For initial inspections, these two types of vehicles are combined in a single queue because initially they share the same equipment (i.e. weighing machine). Queuing areas for light and heavy vehicles at other VIC's are usually adjacent to each other. There are also two separate traffic wardens at each light and heavy queuing area. This shows that two numbering systems, exist for a different type of vehicle.

There are six inspection lanes with three light and three heavy lanes. Currently lanes 1 and 2 have been allocated for finance vehicles, lane 3 for light vehicles and lanes 4 -6 are for heavy vehicles.

Generally, the layout of a VIC affects its inspection flow. At VIC JB, as shown in **Figure 1.1**, we can see that there are four different flows, indicated by four different arrows. Each arrow represents a different type of inspection which flows through the system, throughout the inspection process. Private and accident inspections are loc ated separately from other inspections to speed up the inspection process.

Generally, the sequence of machines in an inspection lane in VIC JB is similar to other VIC's, with a slight difference. **Figure 1.2** shows the sequence of inspection machines located on one of the light inspection lanes and on one of the heavy inspection lanes. Both lanes are preceded by a smoke meter and end with a printer to generate results. In other centers like Wangsa Maju, a printer is shared among other lanes and is located at

one place. When a vehicle has completed its inspection, every driver will go to the printer to obtain the corresponding result.

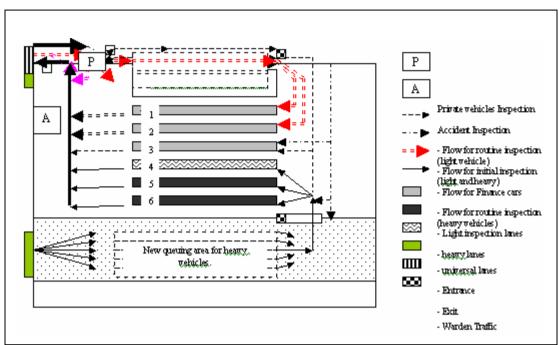


Figure 1.1 : Layout of VIC JB

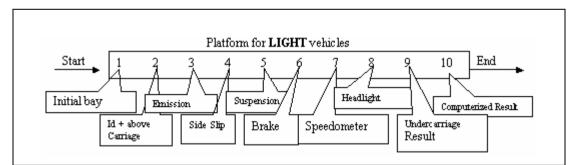


Figure 1.2 : Sequence of Inspection Machines Located on Light Inspection Lane

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1.3.2 Domain Problem – Finance Car Inspection Lane

An inspection flow is different from an inspection procedure as it describes the activity that takes place on the day of inspection. The flow of inspection is slightly different from one VIC to another due to their unique physical layouts. Generally, upon arrival at a VIC, the vehicle owner or driver first shows the security guard related documents as proof of identification of his vehicle and himself. This allows him to enter the VIC premises. Only a vehicle with the correct appointment date and time is permitted to enter the VIC. This is to prevent the premises from becoming over crowded.

The vehicle owner then has to drive the vehicle to a specific queuing area, depending on the required type of inspection. He then has to obtain a queue number from a traffic warden. At the VIC, a numbering system is implemented as there are more than one inspection lane and a sequence of customers keep on coming.

While queuing and waiting to be inspected in the lane, vehicle owner makes a payment as his vehicle approaches the payment counter. When his number is called, his vehicle is passed on to a Vehicle Examiner who then takes the vehicle to the inspection lane and the inspection process starts. After inspection is completed, the result is printed and given to the vehicle owner. Actions regarding the result are then taken by the vehicle owner as described earlier. See **Figure 1.3** for the general procedure in Finance Car Inspection Lane.

The constraints for finance company vehicles are:

- i. Number of lane considered is one.
- ii. No. of inspection considered is ten.
- iii. Rule of inspection is that if all three major items passed, the result is pass.
- iv. Number of major item is three.

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- v. The three major items named chassis number, engine number and original body.
- vi. Maximum number of vehicle examiner is four.

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- vii. Time stagger is one hour (but still not been used in this analysis).
- viii. Time slot is assumed as complete time with waiting time at initial bay.
- ix. Complete time is assumed as processing time with movement period and reserve time.

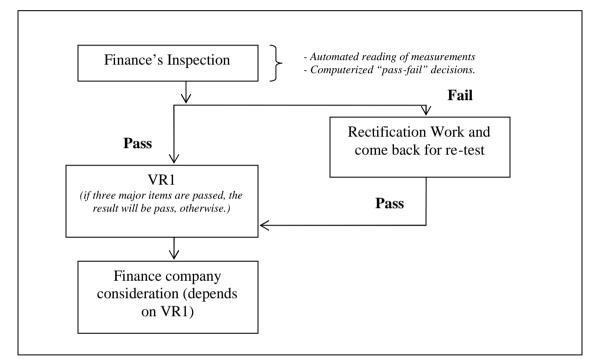


Figure 1.3: General procedure for a finance car inspection

- x. Operation hour in this study is defined from 8 a.m. until 5 p.m.
- xi. There is no tea break defined in this study. It is based on the real operation at the time of data been collected.
- xii. There is no lunch hour defined in this study. It is also based on the real operation at the time of data been collected.
- xiii. Each vehicle must have a valid time slot of appointment to be inspected.But, if there was any vehicle that came at the wrong date and time, it was also been inspected.

1.4 Vehicle Inspection Appointment Schedule

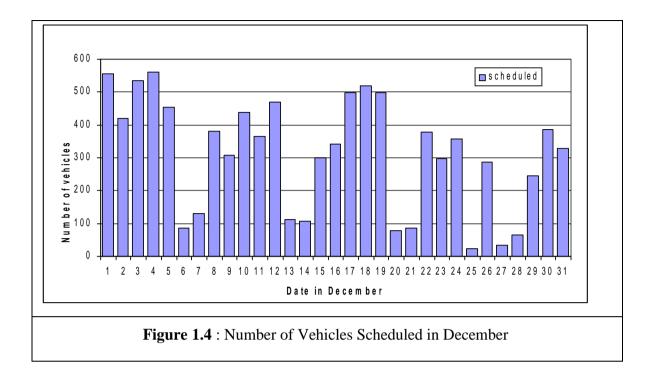
The date and time for a routine inspection is fixed by an appointment system. It is set during the previous routine inspection by a computerized vehicle appointment scheduling system. The system generates a specific inspection date and time for a particular vehicle. A system known as "staggered system" has been applied at the VIC and has significantly reduced congestion at the premises. The staggered system schedules the vehicles' arrival times on an hourly basis. Vehicles should arrive at the allotted times, otherwise they will not be allowed to enter the premises. However the system has been identified as possessing several weaknesses, such as being very rigid, and caters for routine inspections only.

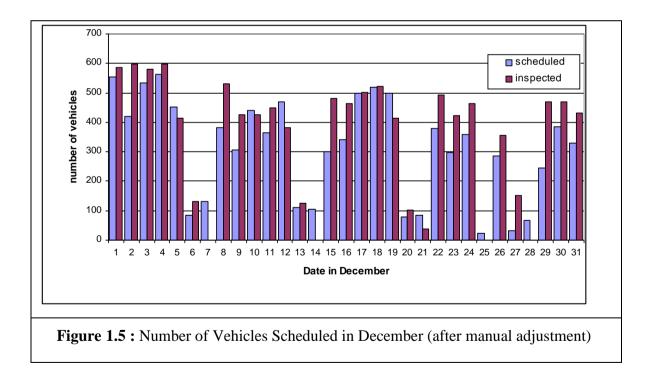
The system does not take into account constraints like holidays. **Figure 1.4** shows the number of vehicles, scheduled for inspection at VIC JB, each day in De cember 2004. The x axis represents the date and note that 1st is on Monday, 2nd is on Tuesday and so on. Sundays fell on 7th, 14th, 21st and 28th December. As noted in figure 6 there are vehicles which were scheduled for inspection on weekends.

To overcome this problem, manual adjustments are made by the Center Manager. **Figure 1.5** shows the numbers of vehicles which have been re-scheduled via "manual adjustment". From the figure, we can see that vehicles previously scheduled for a Sunday have mostly been shifted to the next working day.

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Another factor which makes the system rigid is that it does not consider the maximum capacity of vehicles that can be inspected per hour. The present staggered system sets the next routine inspection time (which is scheduled for exactly six months after the issue of a pass certificate) for a particular vehicle, according to the vehicle's current completion time of inspection. The weaknesses of the vehicle scheduling system at the VIC have some implications on the variability of daily demand, accuracy of vehicles' inspection schedule and workload planning.

Variability in daily demand is created by "manual adjustment" process and "drive -in" vehicles. "Manual adjustment" activity has largely caused the number of vehicles scheduled on specific days to be shifted to the next working days. Hence, the number of vehicles to be inspected on that working day increases. This situation is worsened by "drive-in" vehicles.

When a manual adjustment is done, the altered inspection schedule becomes packed. The situation is worsened by the arrival of drive-in vehicles which is quite high. A schedule which continues to get packed will have smaller inter-arrival times. A smaller inter arrival compared to a constant service time would result in longer waiting times.

Apart from increasing waiting times, the weaknesses of the vehicle scheduling system at the VIC have resulted in difficulties of determine the actual workload for a specific vehicle examiner, to be done beforehand. This has disabled the planning of an effective staff workload which finally creates a situation of "a fixed workforce handlin g variable demand".

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CHAPTER 2

LITERATURE REVIEW

2.1 Distributed Artificial Intelligence (DAI)

Distributed artificial intelligence (DAI) was a subfield of Artificial intelligence research dedicated to the development of distributed solutions for complex problems regarded as requiring intelligence. These days DAI has been largely supplanted by the field of Multi - Agent Systems. See the paper by Inman and Hewitt on some of the limitations of classic DAI.

Background

Main streams in DAI research included the following:

- Parallel problem solving: mainly deals with how classic AI concepts can be modified, so that multiprocessor systems and clusters of computers can be used to speed up calculation.
- Distributed problem solving (DPS): the concept of agen t, autonomous entities that can communicate with each other, was developed to serve as an abstraction for developing DPS systems. See below for further details.

• Multi-Agent Based Simulation (MABS): a branch of DAI that builds the foundation for simulations that need to analyze not only phenomena at macro level but also at micro level, as it is in many social simulation scenarios.

2.2 Complexity of Constraints Satisfaction

The computational complexity of constraint satisfaction has mainly been studied for discriminating between tractable and intractable classes of constraint satisfaction problems on finite domains.

Solving a constraint satisfaction problem on a finite domain is an NP -complete problem in general. Research has shown a number of polynomial -time sub cases, mostly obtained by restricting either the allowed domains or constraints or the way constraints can be placed over the variables. Research has also established relationship of the constraint satisfaction problem with problems in other areas such as finite model theory and databases.

Overview

Establishing whether a constraint satisfaction problem on a finite domain has solutions is an NP complete problem in general. This is an easy consequence of a number of other NP complete problems being expressible as constraint satisfaction problems. Such other problems include propositional satisfiability and three -colorability.

Tractability can be obtained by considering specific classes of constraint satisfaction problems. As an example, if the domain is bina ry and all constraints are binary, establishing satisfiability is a polynomial -time problem because this problem is equivalent to 2-SAT, which is tractable. Research has shown a number of tractable sub cases. Some of these classes are based on restricting the allowed domains or relations, some on restricting the way constraints are placed over variables, and some on both kinds of restrictions. One line of research used a correspondence between constraint satisfaction problem and the problem of establishing the existence of a homomorphism between two relational structures. This correspondence has been used to link constraint satisfaction with topics traditionally related to database theory.

A considered research problem is about the existence of dichotomies a mong sets of restrictions. This is the question of whether a set of restrictions contains only polynomial - time restrictions and NP-complete restrictions. This question is settled for some sets of restrictions, but still open for the set of all restrictions based on a fixed domain and set of relations, as of 2005. This is considered by some authors the most important open question about the complexity of constraint satisfaction.

Restrictions

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Tractable subcases of the general constraint satisfaction problems can be obtained by placing suitable restrictions on the problems. Various kinds of restrictions have been considered.

2.2.1 Structural and relational restrictions

Tractabiliy can be obtained by restricting the possible domains or constraints. In particular, two kinds of restrictions have been considered:

- *relational restrictions* bounds the domain and the values satisfying the constraints;
- *structural restrictions* bounds the way constraints are distributed over the variables.

More precisely, a relational restriction specifies a *constraint language*, which is a domain and a set of relations over this domain. A constraint satisfaction problem meets this restriction if it has exactly this domain and the relation of each constraint is in the given set of relations. In other words, a relational restriction bounds the domain and the set of satisfying values of each constraint, but not how the constraints are placed over variables. This is instead done by structural restrictions. Structural restriction can be checked by looking only at the scopes of constraints (their variables), ignoring their relations (their set of satisfying values).

A constraint language is tractable if there exists a polynomial algorithm solving all problems based on the language, that is, using the domain and relations specified in the domain. An example of a tractable constraint language is that of binary domains and binary constraints. Formally, this restriction corresponds to allowing only domains of size 2 and only constraints whose relation is a binary relation. While the second fact implies that the scopes of the constraints are binary, this is not a structural restriction because it does not forbid any constraint to be placed on an arbitrary pair of variables. Incidentally, the problem becomes NP complete if either restriction is lifted: binary constraints and ternary domains can express the graph coloring problem, while ternary constraints and binary domains can express 3-SAT; these two problems are both NP-complete.

An example of a tractable class defined in terms of a structural restriction is that of binary acyclic problems. Given a constraint satisfaction problem with only binary constraints, its associated graph has a vertex for every variable and an edge for every constraint; two vertices are joined if they are in a constraint. If the graph of a problem is acyclic, the problem is called acyclic as well. The problem of satisfiability on the class of binary ayclic problem is tractable. This is a structural restriction because it does not pl ace any limit to the domain or to the specific values that satisfy a constraints; rather, it restricts the way constraints are placed over variables.

While relational and structural restrictions are the ones mostly used to derive tracable classes of constraint sastisfaction, there are some tractable classes that cannot be defined by relational restrictions only or structural restrictions only. The tractable class defined in terms of row convexity cannot be defined only in terms of the relations or only in t erms of the structure, as row convexity depends both on the relations and on the order of variables (and therefore cannot be checked by looking only at each constraint in turn).

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2.2.2 Uniform and non-uniform restrictions

The subcase obtained by restricting to a finite constraint language is called a *non-uniform problem*. These problems are mostly considered when expressing constraint satisfaction in terms of the homomorphism problem, as explained below. *Uniform problems* were also defined in the settings of homomorphism problems; a uniform problem can be defined as the union of a (possibly infinite) collection of non -uniform problems. A uniform problem made of an infinite set of non -uniform problems can be intractable even if all these non-uniform problems are tractable.

2.2.3 Tree-based restrictions

Some considered restrictions are based on the tractability of the constraint satisfaction problem where the constraints are all binary and form a tree over the variables. This is a structural restriction, as it can be checked by looking only at the scopes of the constraints, ignoring domains and relations.

This restriction is based on primal graph of the problem, which is a graph whose vertices are the variables of the problem and the edges represent the presence of a constraint between two variables. Tractability can however also be obtained by placing the condition of being a tree to the primal graph of problems that are reformulations of the original one.

Equivalence conditions

Constraint satisfaction problems c an be reformulated in terms of other problems, leading to equivalent conditions to tractability. The most used reformulation is that in terms of the homomorphism problem.

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2.3 Constraints satisfaction and the homomorphism problem

A link between constraint satisfaction and database theory has been provided in the form of a correspondence between the problems of constraint satisfiability to the problem of checking whether there exists a homomorphism between two relational structures. A relational structure is a mathematical representation of a database: it is a set of values and a set of relations over these values. Formally, $A = (V, R_1^A, \dots, R_n^A)$, where each R_i^A is a relation over V, that is, a set of tuples of values of V.

A relational structure is different from a constraint satisfaction problem because a constraint is a relation *and* a tuple of variables. Also different is the way in which they are used: for a constraint satisfaction problem, finding a satisfying assignment is the main problem; for a relation structure, the main problem is finding the answer to a query.

The constraint satisfaction problem is however related to the problem of establishing the existence of a homomorphism between two relational structures. A homomorphism is a function from the values of the first relation to the values of the second that, when applied to all values of a relation of the first structure, turns it into a subset of the corresponding relation of the second structure. Formally, *h* is a homomorphism from $A = (V, R_1^A, \dots, R_n^A)$ to $B = (D, R_1^B, \dots, R_n^B)$ if it is a function from *V* to *D* such that, if $(x_1, \dots, x_k) \in R_i^A$ then $(h(x_1), \dots, h(x_k)) \in R_i^B$.

A direct correspondence between the constraint satisfaction problem and the homomorphism problem can be established. For a given constraint satisfaction problem, one can build a pair of relational structures, the first encoding the variables and the signatures of constraints, the second encoding the domains and the relations of the constraints. Satisfiability of the constraint satisfaction problem corresponds to finding a value for every variable such that replacing a value in a signature makes it a tuple in the relation of the constraint. This is possible exactly if this evaluation is a homomorphism between the two relational structures. The inverse correspondence is the opposite one: given two relational structures, one encodes the values of the first in the variables of a constraint satisfaction problem, and the values of the second in the domain of the same p roblem. For every tuple of every relation of the first structure, there is a constraint having as values the correspondent relation of the first structure. This way, a homomorphism corresponds to mapping every scope of every constraint (every tuple of ever y relation of the first structure) into a tuple in the relation of the constraint (a tuple in the corresponding relation of the second structure).

A non-uniform constraint satisfaction problem is a restriction where the second structure of the homomorphism problem is fixed. In other words, every relational structure defines a non-uniform problem, that of telling whether a relation structure is homomorphic to it. A similar restriction can be placed on the first structure; for any fixed first structure, the homomorphism problem is tractable. A uniform constraint satisfaction problem is an arbitrary restriction to the sets of structures for the first and second structure of the homomorphism problem.

2.3.1 Conjunctive query evaluation and containment

Since the homomorphism problem is equivalent to conjunctive query evaluation and conjunctive query containment, these problems below are equivalent to constraint satisfaction as well.

i) Join evaluation

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Every constraint can be viewed as a table in a database, where the variables are interpreted as attributes names and the relation is the set of records in the table. The solutions of a constraint satisfaction problem are the result of an inner join of the tables representing its constraints; therefore, the problem of existence of solutions can be reformulated as the problem of checking whether the result of an inner join of a number of tables is empty.

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Dichotomy theorems

Some constraint languages (or non-uniform problems) are known to be polynomial-time, and some others are known to be NP-complete. However, it is possible that some constraint languages are neither. It is indeed known by Ladner's theorem that, if P is not equal to NP, there exists problems in NP that are neither polynomial -time nor NP-hard. As of 2006, it is not known if such problems include a constraint language. If this is not the case, the set of all constraint languages could be divided exactly in the set of polynomial-time and NP-complete problems, that is, this set has a dichotomy.

Some partial results are however known for some sets of constraint languages. The most known such result is Schaefer's dichotomy theorem, which proves the existence of a dichotomy in the set of constraint languages on a binary domain. More precisely, it proves that a relation restriction on a binary domain is tractable if all its relations belong to one of six classes, and is NP-complete otherwise. Bulatov proved a dichotomy theorem for domains of three elements.

Another dichotomy theorem for a subset of the constraint lan guages is the Hell-Nesetril theorem, which shows a dichotomy for problems on binary constraints having all the same fixed symmetric relation. In terms of the homomorphism problem, every such problem is equivalent to the existence of an homomorphism from a relational structure to a given fixed undirected graph (an undirected graph is a relational structure with a single binary symmetric relation). The Hell-Nesetril theorem proves that every such problem is either polynomial-time or NP-complete. More precisely, the problem is polynomial-time if the graph is 2-colorable, that is, it is bipartite, and is NP-complete otherwise.

Sufficient conditions for tractability

Some complexity results prove that some restrictions are polynomial without giving proving that all other possible restrictions of the same kind are NP -hard. A sufficient condition for tractability is related to expressible in Datalog. A Boolean Datalog query

gives a truth value to a set of literals over a given alphabet, each literal being an expression of the form $L(a_1,...,a_n)$; as a result, a Boolean Datalog query expresses a set of sets of literals, as it can be considered semantically equivalent to the set of all sets of literals that it evaluates to true.

On the other hand, a non-uniform problem can be seen as a way for expressing a similar set. For a given non-uniform problem, the set of relations that can be used in constraints is fixed; as a result, one can give unique names R_1, \ldots, R_n to them. An instance of this non-uniform problem can be then written as a set of literals of the form $R_i(x_1, \ldots, x_n)$. Among these instances/sets of literals, some are satisfiable and some are not; whether a set of literals is satisfiable depends on the relations, which are specified by the non-uniform problem. In the other way around, a non-uniform problem tells which sets of literals represent satisfiable instances and which ones represent unsatisfiable instances. Once relations are named, a non-uniform problem expresses a set of set s of literals: those associated to satisfiable (or unsatisfiable) instances.

A sufficient condition of tractability is that a non -uniform problem is tractable if the set of its unsatisfiable instances can be expressed by a Boolean Datalog query. In other words, if the set of sets of literals that represent unsatisfiable instances of the non uniform problem is also the set of sets of literals that satisfy a Boolean Datalog query, then the non-uniform problem is tractable.

ii) Local consistency

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Satisfiability can sometimes be established by enforcing a form of local consistency and then checking the existence of an empty domain or constraint relation. This is in general a correct but incomplete unsatisfiability algorithm: a problem may be unsatisfiable even i f no empty domain or constraint relation is produced. For some forms of local consistency, this algorithm may also require exponential time. However, for some problems and for some kinds of local consistency, it is correct and polynomial -time.

The following conditions exploit the primal graph of the problem, which has a vertex for each variable and an edge between two nodes if the corresponding variables are in a constraint. The following are conditions on binary constraint satisfaction problems where enforcing local consistency is tractable and allows establishing satisfiability:

- 1. enforcing arc consistency, if the primal graph is acyclic
- 2. enforcing directional arc consistency for an ordering of the variables that makes the ordered graph of constraint having width 1 (such an ordering exists if and only if the primal graph is a tree, but not all orderings of a tree generate width 1)
- 3. enforcing strong directional path consistency for an ordering of the variables that makes the primal graph having induced width 2.

A condition that extends the last one holds for non -binary constraint satisfaction problems. Namely, for all problems for which there exists an ordering that makes the primal graph having induced width bounded by a constant i, enforcing strong directional i-consistency is tractable and allows establishing satisfiability.

iii) Tree-based conditions

Constraint satisfaction problems composed of binary constraints only can be viewed as graphs, where the vertices are variables and the edges represent the presen ce of a constraint between two variables. This graph is called the Gaifman graph or primal graph of the problem.

If the primal graph of a problem is acyclic, establishing satisfiability of the problem is a tractable problem. This is a structural restriction, as it can be checked by looking only at the scopes of the constraints, disregarding their relations and the domain. An acyclic graph is a forest, but connectedness is usually assumed; as a result, the condition that is usually considered is that primal graphs are trees.

This property of tree-like constraint satisfaction problems is exploited by decomposition methods, which convert problems into equivalent ones that only contain binary constraints arranged as a tree. The variables of these problems corres pond to sets of variables of the original problem; the domain of such a variable is obtained by considering some of the constraints of the original problem whose scope is contained in the corresponding original set of variables; constraints of these new pr oblems represent equality of variables that are contained in two sets.

If the graph of one such equivalent problem is a tree, the problem can be solved efficiently. On the other hand, producing one such equivalent problem may be not be efficient because of two factors: the need to determine the combined effects of a group of constraints over a set of variables, and the need to store all tuples of values that satisfy a given group of constraints.

Necessary condition for tractability

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A necessary condition for the tractability of a constraint language based on the universal gadget has been proved. The universal gadget is a particular constraint satisfaction problem that was initially defined for the sake of expressing new relations by projection.

2.4 Distributed Constraints Satisfaction Problem (DisCSP)

DisCSP is a framework for describing a problem in terms of constraints that are known and enforced by distinct participants (agents). The framework was used under different names in the 1980's. The first known usage with the current name is in 1990. The constraints are described on some variables with predefined domains, and have to be assigned to the same values by the different agents. Problems defined with this framework can be solved by any of the algorithms t hat are proposed for it.

The currently most well known paper is the one which introduced in 1992 as asynchronous complete algorithm for solving such problems is Asynchronous backtracking (ABT). A large number of such techniques are now available, See **Table 2.1** below.

Year	Technique	Abbreviation	Characteristics
1992	Asynchronous Backtracking	ABT	Static ordering, complete search
1994	Asynchronous Weak- Commitment	AWC	Reordering, fast, complete search
1995	Distributed Breakout Algorithm	DBA	Incomplete but fast
2000	Distributed Forward Chaining	DFC	Slow, comparable to ABT
2000	Asynchronous Aggregation Search	AAS	Aggregation of values in ABT
2001	Maintaining Asynchronously Consistencies	DMAC	The fastest algorithm
2001	Asynchronous Backtracking with Reordering	ABTR	Reordering in ABT with bounded nogoods
2002	Secure Computation with Semi-Trusted Servers	-	Security increases with the number of trustworthy servers
2003	Secure Multiparty Computation for Solving DisCSPs	MPC- DisCSP1-	Secure if half (1/2) of the participants are trustworthy.
		MPC- DisCSP4	

Table 2.1: Distributed Constraints Satisfaction (DisCS) Techniques

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2.5 Agent Technology

What is Agent

Agent is anyone who acts on behalf or in the interest of somebody else. So most agent are designed to act as or for the user to help execute some task or operation. Agents come in various form for various purpose. Here, an agent can be a human being, a robot, or a computer program. Consequently, a more precise and very general definition can be provided for agents is it can do routine work for users or assist them with complicated and finished their tasks.

In the internet definition, it defines that an agent (also called an intelligent agent and sometime called a bot) is a program that gathers information or performs some other service without your immediate presence and on some standard schedule. Typically, an agent is an entity that is able to carry out some task usually to help human user. Agent can be biologic, robotic or computational. In general we can describe the agent as referring to a component of software or hardware which is able of acting in order to accomplish tasks on behalf of its user.

<u>Software Agent</u>

A software agent is a computer program (more to software) that can accept tasks from its human user can figure out which actions to perform in order to solve these tasks and can actually perform these task actions without user aided.

2.5.1 Agent Properties

The behavior of the agent must have at least the following properties: -

Autonomy: Agent can operate on their own without the need for human guidance and it also called as proactive because these agents can take the initiative and make decision independently related to control; although an agent may interact with its environment, the processes performed by an agent are in full control of the agent itself.

Reactivity: leading the agent to rapidly respond to new information perceived from its environment. Reactivity indicates that the agent will respond to perceived changes in the environment and able to handle changing environment.

Proactively: agents should not simply act in response to their environment; they should be able to exhibit opportunistic, goal-directed behavior and take their initiative where appropriate. In situation where the environment changes or the goal achievement fails, the agent will still to find way to reach the goal

Learning: React and interact with their external environment, from observ ation or by user's instructions.

Flexible: Agents have flexible behavior, because they can achieve in several ways. To reach a goal an agent chooses a plan to follow. If the plan fails, the agent can choose between other possibilities to obtain the desire d result.

Social ability / Cooperation: is the central property in *multi*-agent systems. Agents must be able to communicate and cooperate with other agents and humans to complete their own problem solving and help others with their activities. It also called as cooperation, which is the ability to interact with other agent, cooperate with each others and interaction with possible human via some communication language.

Other non-mandatory properties of agents include *adaptability*, *mobility*, and *intelligence*. Although the term software agent is sometimes used as a synonym for intelligent software systems, software agents do not have to be intelligent. In software engineering, for instance, the ability of an agent to communicate and cooperate with other systems in a flexible manner, and the ability of a mobile agent to migrate to another

computer providing resources via suitable network links are considered more fundamental than any form of intelligence.

2.5.2 How they applied this technology in a system

Agents propose new possibilities and solution to answer the problem suitable to their properties and characteristics. Some of the characteristic of agent is know the preferences of their user. The efficiency of the scheduling process used and the quality of the schedule generated, then affects the working of an organization to a large extends. The used of agent can developed more efficient schedule because it can save time and effort on the part of human. There is also a communication between agent and user. The age nts manage the scheduling process on behalf the individual user, without human interaction. Each agent can access to calendar and preference information of its user.

Generally, agent has initial information about schedule the meeting. That means, the agents have to learn and get to know the habit and preference of user by observing the action of the user. Some of the actions are marking suggestion, collecting or receive feedback either positive or negative and the last is suggestion from the other user. So from this way, the user comes to trust the agent and decide to delegate their meeting scheduling decision. In this approach, agent also provide autonomous scheduling capabilities between intelligent agent, where these agent use knowledge about the priorities of people to perform their task such as automatically screening, directing and even responding to information.

The meeting to be schedule must have a day, a start time and place, and the problem is solve when all agent reach an agreement on values for these attributes. Agents negotiate by having one agent propose a meeting. The agent who initiates the new meeting called as host agent. While the meeting is start, the other attendee or invitee agent have ability to make decision on behalf of their user to respond the invitation by accepting or rejecting them, based on weather or not it fit their own user's schedules. User will be supervising a

group of agent such as in their activity and monitoring for any change in the schedule. Then agent predicts a response and negotiates with other agent to suggest an optimal time for the meeting.

2.6 Intelligent Agent

The terms "agent" and "intelligent agent" are ambiguous and have been used in two different, but related senses, which are often confused.

In computer science, an intelligent agent (IA) is a software agent that assists users and will act on their behalf, in performing non-repetitive computer-related tasks. An agent in the sense of the word is like an insurance agent or travel agent. While the working of software agents used for operator assistance or data mining (sometimes referred to as bots) is often based on fixed pre-programmed rules, "intelligent" in this context is often taken to imply the ability to adapt and learn.

In artificial intelligence, the same term is used for intelligent actors, which observe and act upon an environment, to distinguish them from intelligent thinkers isolated from the world. An agent in this sense of the word is an entity that is capable of perception and action. Such an agent might be a robot or an embedded real time software system - and is intelligent if it interacts with its environment in a manner that would normally be regarded as intelligent if that interaction were carried out by a human being.

There is no reason why these two notions of intelligent agent need to be related. An intelligent agent in the first sense might be implemented using conventional software techniques and display no more intelligence than a conventional computer program. On the other hand, an intelligent agent in the second sense might be wholly autonomous, carrying out its own agenda, and acting as an agent for no one.

In the artificial intelligence sense of the term, there are multiple types of agents and sub - agents. For example:

- 1. Physical Agents A physical agent is an entity which *percepts* through sensors and *acts* through actuators.
- Temporal Agents A temporal agent may use time based stored information to offer instructions or data *acts* to a computer program or human being and takes program inputs *percepts* to adjust its next behaviors.

A simple agent program can be defined mathematically as an agent function which maps every possible percepts sequence to a possible action the agent can perform or to a coefficient, feedback element, function or c onstant that affects eventual actions:

$$f: P * - > A$$

The program agent, instead, maps every possible percept to an action.

It is possible to group agents into five classes based on their degree of perceived intelligence and capability:

- 1. simple reflex agents;
- 2. model-based reflex agents;
- 3. goal-based agents;
- 4. utility-based agents;
- 5. learning agents.

1. Simple reflex agents

Simple reflex agents act only on the basis of the current percept. The agent function is based on the condition-action rule:

if condition then action rule

This agent function only succeeds when the environment is fully observable. Some reflex agents can also contain information on their current state which allows them to disregard conditions whose actuators are already triggered.

2. Model-based reflex agents

Model-based agents can handle partially observable environments. Its current state is stored inside the agent maintaining some kind of structure which describes the part of the world which cannot be seen. This behavior requires information on how the world behaves and works. This additional information completes the "World View" model.

3. Goal-based agents

Goal-based agents are model-based agents which store information regarding situations those are desirable. This allows the agent a way to choo se among multiple possibilities, selecting the one which reaches a goal state.

4. Utility-based agents

Goal-based agents only distinguish between goal states and non -goal states. It is possible to define a measure of how desirable a particular state is. This measure can be obtained through the use of a *utility function* which maps a state to a measure of the utility of the state.

5. Learning agents

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In some literature IAs are also referred to as *autonomous intelligent agents*, which mean they act independently, and will learn and adapt to changing circumstances. IA systems should exhibit the following characteristics:

- learn and improve through interaction with the environment (embodiment)
- adapt online and in real time
- learn quickly from large amounts of data
- accommodate new problem solving rules incrementally
- have memory based exemplar storage and retrieval capacities
- have parameters to represent short and long term memory, age, forgetting, etc.
- be able to analyze itself in terms of behavior, error and success.

To actively perform their functions, Intelligent Agents today are normally gathered in a hierarchical structure containing many "sub-agents". Intelligent sub-agents process and perform lower level functions. Taken together, the intelligent agent and sub -agents create a compete system that can accomplish difficult tasks or goals with behaviors and responses that display a form of intelligence.

Some of the sub-agents (not already mentioned in this treatment) that may be a part of an Intelligent Agent or a complete Intelligent Agent in them are:

- 1. Temporal Agents (for time-based decisions);
- 2. Spatial Agents (that relate to the physical real -world);
- Input Agents (that process and make sense of sensor inputs example neural network based agents neural network);
- 4. Processing Agents (that solve a problem like speech recognition);
- 5. Decision Agents (that are geared to decision making);
- Learning Agents (for building up the data structures and database of other Intelligent agents);
- 7. World Agents (that incorporate a combination of all the other classes of agents to allow autonomous behaviors).

A very limited set of agents, that might be classified as semi -intelligent due to their lack of complexity, decision making, extremely limited world view and learning capacity can be found in the reference: Third Canadian Edition of "Management Information Systems for the Information Age". This document suggests that there are only four essential types of Intelligent Agents:

- 1. Buyer agents or shopping bots
- 2. User or personal agents
- 3. Monitoring-and-surveillance agents
- 4. Data Mining agents

Buyer Agent

Buyer agents travel around a network (i.e. the internet) retrieving information about goods and services. These agents, also known as 'shopping bots', work very efficiently for commodity products such as CDs, books, electronic components, and other one -size-fits-all products. *Amazon.com* is a good example of a shopping bot. The website will offer you a list of books that you might like to buy on the basis of what you're buying now and what you have bought in the past.

User or Personal Agents

User agents, or personal agents, are intelligent agents that take action on your behalf. In this category belong those intelligent agents that already perform, or will shortly perform, the following tasks:

- Check your e-mail, sort it according to priority (your priority), and alert you when good stuff comes through like college acceptance letters
- Play computer games as your opponent or patrol game areas for you
- Assemble customized news reports for you. There are sev eral versions of these, CNN being a prime example
- Find information for you on the subject of your choice
- Fill out forms on the Web automatically for you, storing your information for future reference
- Scan Web pages looking for and highlighting text that constitutes the "important" part of the information there
- "Discuss" topics with you ranging from your deepest fears to sports

Monitoring and Surveillance Agents

These agents, also known as "predictive agents", are intelligent agents that observe and report on equipment. For example, NASA's Jet Propulsion Laboratory has an agent that monitors inventory, planning, and scheduling equipment ordering to keep costs down, as well as food storage facilities. These agents usually monitor complex computer networks that can keep track of the configuration of each computer connected to the network.

Data Mining Agents

A data mining agent operates in a data warehouse discovering information. A 'data warehouse' brings together information from lots of different sources. 'Data mining' is the process of looking through the data warehouse to find information that you can use to take action, such as ways to increase sales or keep customers who are considering defecting. 'Classification' is one of the most common types of data mining, which finds patterns in information and categorizes them into different classes. Data mining agents can also detect major shifts in trends or a key indicator and can detect the presence of new information and alert you to it.

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CHAPTER 3

METHODOLOGY

3.1 Introduction

Some organization especially in safety-critical area requires that certain methodologies be used as it increases the quality of their software and also offer a certain degree of legal protection to all parties involve. Methodology helps to manage with the complexity of the software development process. Methodology may provide a framework for application of technique and resources at appropriate times during the development process. According to Bernard a good software methodology is a methodology which:-

> "can be describes qualitatively, as well as qualitatively, can be used repeatedly, each time achieving similar results, can be taught to others within a reasonable timeframe, can be applied by others with a reasonable level of success"

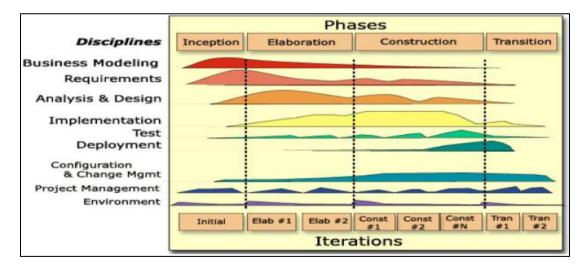
This chapter emphasizes on system development process which focuses on project methodology, a brief explanation on Unified Modeling Language as a technique for design the functionality of system and the end of this section this chapter will be concluded by the summary. In this area the RUP methodology has been used as a technique and guidelines for software development process. There are including the workflow on RUP and also elaborate in detail on their phase such as inception, elaboration, construction and testing phase.

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3.2 Software Development Process

A software development process is describing software life cycle and software process. It begins with architecture, moves through analysis, design, and implementation, undergoes testing and finally is delivered to user. It also covers how to develop the software in the best way, a guideline and manner to develop the best quality of product by variety of tasks or activities for each phase in project methodology.

There are various difference software development methodologies available for modeling software system. The Rational Unified Process (RUP) has been selected in this project methodology. The choice of RUP methodology for a software development process because this technique is the best practice that can be tailored with large and suitable for all range of project and it aimed at meeting the needs of end users. RUP is an Object Oriented framework which provides guidelines and templates and tools to support software d evelopment process.



Figures 3.1: RUP methodology

Based on the figure above, that is show the RUP software development process which consists four chronological phases. There are Inception, Elaboration, Construction and Transition. The vertical line represents the workflow in RUP and the detail description is shown as below:

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1. Business modeling

Understanding the structure and operation of the business

How the product fits into their organization

2. Requirements

Understand what user need from a product and to define these needs is a way that can be the basis for designing and planning product's content.

3. Analysis and design

Translate requirement into a design for the system

4. Implementation

Construct a working version of the system.

5. Test

Verify that the system meets requirement and minimize errors

6. Deployment

Distribute the product to users.

The Benefits of RUP

RUP provides the system development team with the advantages of RUP best practice s while providing a framework for addressing overall system issues. Some of the benefits of RUP include:

- Provides for ongoing collaboration among business analysts, architects, system engineers, software developers, hardware developers, and testers.
- Provides views that enable teams to address system quality issues in an architecture driven process.
- Provides UML support for systems architecture.
- Scales upward to very large systems.
- Provides workflows for determining hardware and software components.
- Supports concurrent design, and iterative development of hardware and software components.

- Easy to understand and facilitate resource planning and work structuring
- RUP delivers a well-structured framework, divided into phase and workflow, allowing easy navigation within the framework

Phases in RUP Methodology

According to RUP methodology, the phase in RUP had been applied as a guideline to accomplish the development of new reservation scheduling system. Following of that, the detail of description about each phase in RUP methodology and some related works that have been done show as below:-

Inception Phase

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Inception phase aimed to establish a business case. In general, the problem is first defined and analysis of the requirement of the system will be identified by discussion with the user. Included in this phase, the title was proposed, the background of problem in case study has been studied. The scope and objectives for development of this Appointment Scheduling for Reservation System (ASRS) system had been determined and clearly stated. This phase is very important to identify the business needs from the system and also to get the business value from the proposed system.

Following this, the real problem and process had been conducted by observation and int erview session. The real information of their operations has been collected by visited at some of vehicle inspection center. From this study some of the problem was define and the requirement from the client was generated. In order to have better performan ce on inspection appointment, they required for modeling and designing new model system to solve all problems occurred in the manual system especially to reduce the amount of time processing. The literature review was also conducted in this phase. It is covers on technique and a new technology that would be applied in this new system such as Agent Technology, Constraint Logic Programming and others.

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Elaboration Phase

Elaboration is the initial series of iteration which the team does serious investigation, implements the program, build the system architectures, analysis the user requirement specification and find out some high-risk issues.

During elaboration phase, the architecture of new system will be design. In this phase the real data for processing time in each process of vehicle inspection was collected. Data collection had been carried in vehicle inspection center. The processing time was recorded in seconds due to precision. From data collection, the data analysis is initiated by extracting data int o the algorithm. Data analyses is some of steps for transforming data with the aim to generate some algorithm that will serve as an engine to this new system to solve time constraint problem.

Beside of that, the functional requirement was defined requested by user for new reservation scheduling system. The user requirements have been described in Software Requirement Specification (SRS) document. In this document, that is not only cover on user requirement but also the explanation about the performance, reliability, security, maintainability and other critical issues. Finally, the initial design and architecture modeling will be generated using UML design for the proposed system development. The preliminary designs have been developed including use case di agram, class diagram, sequence diagram and activity diagram.

Construction Phase

A coding design may be the most obvious and difficult part on development of system. The system will de coded by programming language. In this process, Java programming lan guage has been selected as a programming language to build ASRS system. During this phase, the unit testing must be conducted together to make sure the functionality of module process was successfully worked and meets the user's requirement.

Transition Phase

After the code is appropriately tested and approved, it is moved into transition phase to ensure the software is available to end user. This phase includes final testing of product and making some minor adjustment based on error feedback. As a part of this phase, the application will be tested to end user and they will evaluate the quality of system. It is very important to have training classes for most users to build excitement into adopting the new system. At the end of this phase is documenting the internal design of software for the purpose of future maintenance and enhancement. Documentation is important as a user guideline to use the new system. The several project documents were produced such as user manual, a plan to support the user and a plan for upgrading the information system in the future. Then all requirement of the system will be installed and the related document will be delivered to end user.

3.3 Software Development Tools

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The well-known Unified Modeling Language (UML) has been c hosen as tools for software development. The UML is a very important part of developing objects oriented software and the software development process. The UML represent a collection of best engineering practices that have proven successful in the modeling of large and complex system. The UML can help even non-programmer to understanding the overall functionality of the system in a way that meets the entire requirement in the system. The modeling design patterns with UML have several advantages include:-

- It quite naturally to object oriented modeling
- Process models can be communicated more easily in a large number of people
- UML provide a large set of diagram which can be used to define both structured and behavior of dynamic software process
- Object oriented modeling support the earlier phase of process model development
- By UML process model it easy to map them to programmed systems.

Phase	Aim	Activity	Previous Works
Inception	To establish a	-define a case study	-analyze the real specification in case
	business case	-investigate the real	study
		problem	-doing some literature review in
		-Make some literature	DisCSP, Agent Technolology and CLP
		review	that would be applied in this system
Elaboration	Analysis System	-data collection	-Collect real data in vehicle inspection
	design and	-data analysis	center
	system	-Initial design	-Analyze the data
	architecture		-Doing some data analysis and extract
			the data to mathematical model
			-Designing the mathematical model
			-Designing system architecture and
			initial design for new system
Construction	Code writing	-Full design	-Designing full design using UML
	and unit testing	-code writing	diagram for new system
		-unit testing	-Full coding for ASRS system using
			Java programming using Eclipse
			-Unit testing has been conducted in each
			module.
Transition	Documentation	-final testing	-some adjustment has been conducted
	report and	-make some minor	The full documentation has been done
	deliver to end	adjustment	User manual has been done
	user	-user manual	
		-documentation for user	
		-installation new system	

Table 3.1: Summary of RUP phases

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3.3.1 UML diagrams

UML defines various types of diagrams: use case diagram, class diagram, sequence diagram, activity diagram, deployment diagram, component diagram, collaboration diagram and object diagram. The UML diagrams were recommended used as input to test cases. Use case were input functional test specification, sequence diagram were input to integration testing and class diagram were input to unit testing [4]. For further detail, here some description about UML diagram and their function.

<u>Use case Diagram</u>

Use case describes how external actors interact with the software system. Use case diagram consist of name pieces of functionality, the person invoking the functionality which called as actor and possibly the element responsible for implementing the use case. During this interaction an actor generates events to a system. That is a way to capture the functionality and requirement of system. Generally, use cases express the behavior of a system, model the functionality requirement of system using actors and use cases without the excessive detail which often confuses people with a less technical background.

<u>Class Diagram</u>

This diagram shows the structure and behavior of program system. In class diagram contents using design elements such as classes, packages and objects, name, signature, and properties. There are several types of relationship in class diagrams such as association, dependency, generalization, aggregation, composition and realization. E ach relationship is represented in the diagram by a difference type of arrow.

Activity Diagram

That is used to model workflow or business processes and internal operation. It illustrates the dynamic nature of a system by modeling the flow of control f rom activity to activity. An activity represents an operation on some class in the system that results in a change in the state of the system.

Sequence diagram

The sequence diagram will be used to illustrate the design by represent how object interact with each other via message in the operation to fulfill work. This diagram used in elaborating and detailing the behavior of a system. This diagram show the interaction between the object involved in the use case, the actors and flows of events. On horizont al axis it shows of the object that it represented, while on vertical axis it shows the flow or sequence of the event.

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CHAPTER 4

EXPERIMENTAL DATA AND ANALYSIS

4.1 Introduction

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Data Analysis is some sort of steps of transforming data with the aim of extracting the data that we get through particular data collection. In our study, we use both quantitative and qualitative data as our input to be used in the next steps of analysis. Based on analysis of the history of science, Kuhn concludes that :

"Large amounts of qualitative work have usually been prerequisite to fruitful quantification in the physical sciences."

The result of these steps is a number, or a series of numbers that would be used to the next steps of data analysis, or it can be concluded as a result of the study. In our case, we got the results in a form of table, and then we use the figure from the table to be analyzed into the next experimental steps. Generally, results of quantitative methods often presented in tables, charts, graphs or other forms of statistics. Statistics is a mathematical science relating to the collection of data, analysis of the data, and also interpretation or explanation of the data, as well as presentation of data as a final stage.

Marshall et. al in his study outlined the four broad questions that should be considered in the data collection.

- i. In what ways does it help our study?
- ii. What are the advantages or disadvantages of the data?
- iii. What are the contributions will we get from the data?
- iv. How it could be changed or improved?

In general, there are two statistical methods those can be used to collect and analyze, interpret, explain, and present the data. The one which can be used to summarize or describe a collection of data is known as descriptive statistics, while inferential statistics is a method to model the data in a way that accounts for randomness and uncertainty in the observation, and then used to draw inferences about the process or population being studied. Statistical inference or also called as statistical induction comprises the use of statistics to make inferences concerning some unknown aspect of a population. It is distinguished from descriptive statistics.

We use Ms Excel 2003 as an analytical tool to calculate the mean and average values of the data. Then, we use this software to analyze the rest of steps involve in this analyses.

4.2 Data Collection

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Before collecting real data from the case study, we first planed and defined the field and information which relevant for this study. From the observation, we made our strategy to collect only the relevant data which is include the processing time for every point of the vehicle inspection. There are 10 points of vehicle inspection. Every point of inspection is defined by a prefix P1, P2, ..., P10. The details for point of inspection described as follow:

- P1 : Identification check verifies the engine and chassis number of the vehicle to be tested.
- P2 : Above carriage check the vehicle is then scrutinized for surface defects and non-compliance with road transport act 1987.

- 3) P3 : Emission test density of particles in diesel engines and volume of gases
 e.g. HC, CO etc in petrol engines are gauged through free acceleration tests.
- P4 : Side-slip test lateral movements of the vehicle are tested for optimum road-handling criteria.
- 5) P5 : Suspension test checks the effectiveness of the suspension system (springs, shock absorber, and joints) on each axle.
- P6 : Brake test the various brake efficiencies, such as service brake efficiency, dynamic unbalance, run-out, residual force and parking brake are evaluated for maximum performance.
- P7 : Speedometer checks the road speed at the wheels against the value indicated on the vehicle's speedometer.
- 8) P8 : Headlight test measures the intensity and projection of the vehicle's headlights
- 9) P9 : Undercarriage check a vehicle's undercarriage is thoroughly scrutinized for defects while it is tested on an axle play detector.
- 10) P10 : Computerized analysis a state-of-the-art central computer controls and monitors each test phase. All relevant test data is then analyzed and the final result printed out.

Data are collected from various sources as follows:

- i. Observation on environment, management and services.
- ii. Observation on the process of inspections.
- iii. Recording time taken on every point of inspections' process.
- iv. Recording any significant or weird action or event happened during observation.
- v. Request certain appropriate form and information from the organization.

Here are the descriptions of data that we have collected for this study.

- Plat number of the vehicle.
- Brand and model of the vehicle.
- Traveling time to the initial bay.

- Length of time in the initial bay.
- Traveling time to the first point, P1.
- Length of time in the first point, P1.
- Traveling time to the second point, P2.
- Length of time in the second point, P2.
- Traveling time to the third point, P3.
- Length of time in the third point, P3.
- Traveling time to the forth point, P4.
- Length of time in the forth point, P4.
- Traveling time to the fifth point, P5.
- Length of time in the fifth point, P5.
- Traveling time to the sixth point, P6.
- Length of time in the sixth point, P6.
- Traveling time to the seventh point, P7.
- Length of time in the seventh point, P7.
- Traveling time to the eighth point, P8.
- Length of time in the eighth point, P8.
- Traveling time to the ninth point, P9.
- Length of time in the ninth point, P9.
- Traveling time to the tenth point, P10.
- Length of time in the tenth point, P10.
- Traveling time for the Vehicle Examiner (VE) to make u -turn; which is VE has to turn back from the P10 of the completed inspection to the P1 of the new inspection.

Based from the meeting with Quality Assurance Manager, the average processing time for every vehicle being inspected is around five minutes. This figure is based on the production volume of the day which recording the start and end time for every vehicle inspection. **Figure 4.1** below shows the sample of the data that are collected at the vehicle inspection centre.

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Figure 4.1: Sample of Data Collection

4.3 Data Filtering

From the original raw data, data were improved by a little adjustment. Data with major difference will be deleted from the record and the rest of the data will be use for the next step of analysis. There are several ways the data are filtered. In this study the data filtering is done as follows:

i. Sorting the value for each column of point of inspection, one by one, to be in ascending order.

- ii. Then, recognizing if there is any weird value from the first and last part of the overall value in each column. If the value is out of normal range, then the value is marked as a weird value.
- iii. If there is a weird value, it is tried to be rectified. If cannot, deleting a record at all.
- iv. Repeat the steps above to each column in database.

Figure 4.2 below shows the result after the filtering technique is done.

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Figure 4.2 : Sample of datasets after filtering

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4.4 Data Experimental and Analysis

After data filtering, we then compute some descriptive value to get the pattern of the data. The main goal in this step is to obtain the average value from every point of inspection. The values computed are listed below, and shown in **Figure 4.3** and **Figure 4.4**.

- i. Minimum value
- ii. Maximum value
- iii. Number of data
- iv. Average value
- v. Mod value
- vi. Standard deviation

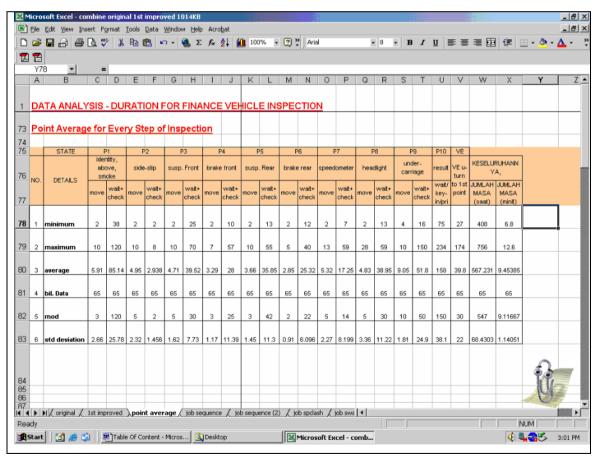
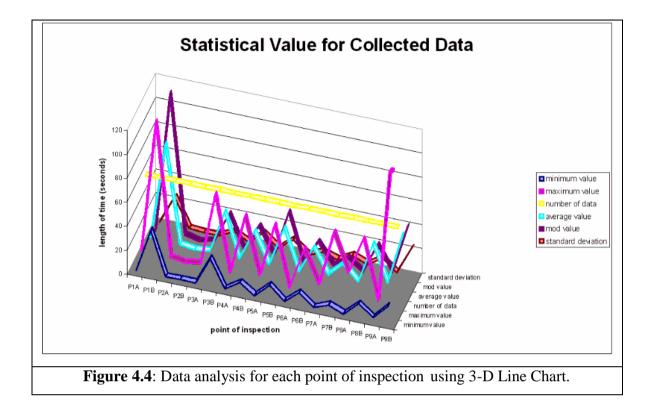


Figure 4.3: Data analysis for each point of inspection.

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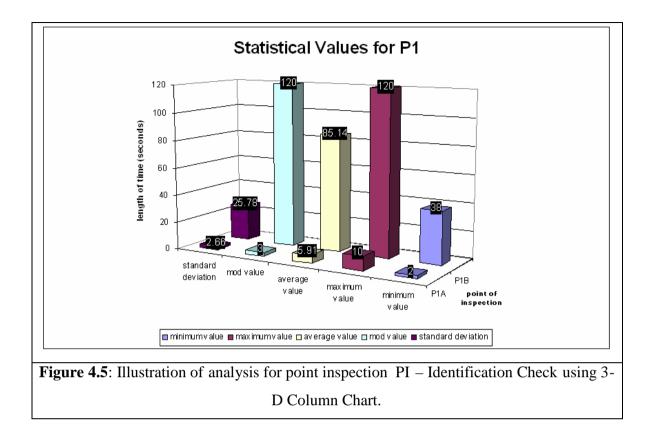


P1 : Identification check

The analysis for this point of inspection is described in **Table 4.1** below, while the illustration of the figures is shown in **Figure 4.5**.

	P1A (traveling time for P1)	P1B (service time for P1)
Minimum value	2 seconds	38 seconds
Maximum value	10 seconds	120 seconds
Number of data	65 vehicles	65 vehicles
Average value	5.91 seconds	85.14 seconds
Mod value	3 seconds	120 seconds
Standard deviation	2.66 seconds	25.78 seconds
Table 4.1: A	Analysis for point inspection PI –	Identification Check

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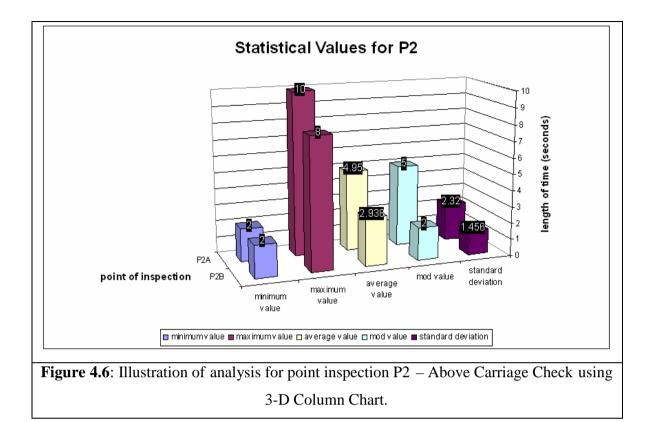
<u>P2</u> : Above carriage check

The analysis for this point of inspection is described in **Table 4.2** below, while the illustration of the figures is shown in **Figure 4.6**.

	P1A (traveling time for P1)	P1B (service time for P1)
Minimum value	2 seconds	2 seconds
Maximum value	10 seconds	8 seconds
Number of data	65 vehicles	65 vehicles
Average value	4.95 seconds	2.938 seconds
Mod value	5 seconds	2 seconds
Standard deviation	2.32 seconds	1.456 seconds
Table 4.2: A	nalysis for point inspection P2 – A	Above Carriage Check

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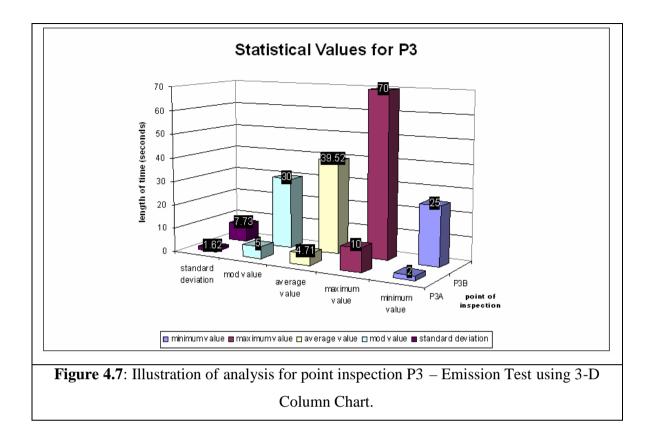


P3 : Emission Test

The analysis for this point of inspection is described in **Table 4.3** below, while the illustration of the figures is shown in **Figure 4.7**.

	P1A (traveling time for P1)	P1B (service time for P1)
Minimum value	2 seconds	25 seconds
Maximum value	10 seconds	70 seconds
Number of data	65 vehicles	65 vehicles
Average value	4.71 seconds	39.52 seconds
Mod value	5 seconds	30 seconds
Standard deviation	1.62 seconds	7.73 seconds
Table 4.	3: Analysis for point inspection P	3 – Emission Test

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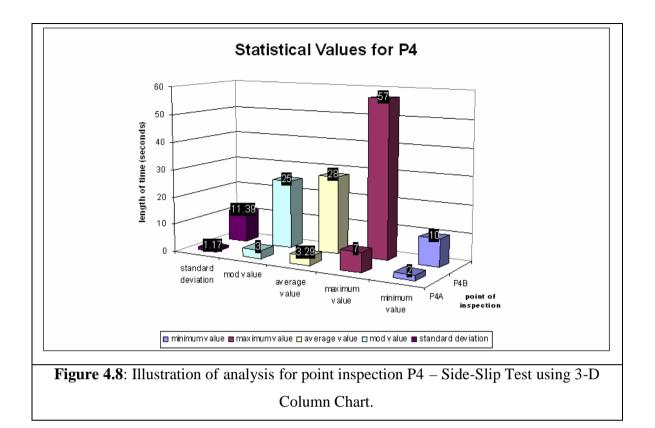
<u>P4 : Side-Slip Test</u>

The analysis for this point of inspection is described in **Table 4.4** below, while the illustration of the figures is shown in **Figure 4.8**.

	P1A (traveling time for P1)	P1B (service time for P1)
Minimum value	2 seconds	10 seconds
Maximum value	7 seconds	57 seconds
Number of data	65 vehicles	65 vehicles
Average value	3.29 seconds	28 seconds
Mod value	3 seconds	25 seconds
Standard deviation	1.17 seconds	11.39 seconds
Table 4.4	Analysis for point inspection P	4 – Side-Slip Test

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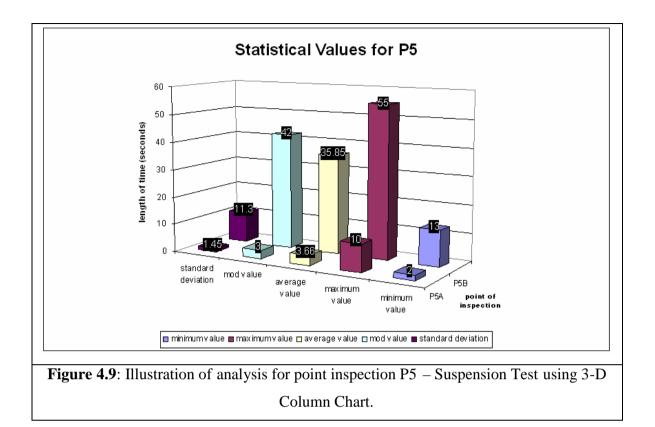
<u>P5</u> : Suspension Test

The analysis for this point of inspection is described in **Table 4.5** below, while the illustration of the figures is shown in **Figure 4.9**.

	P1A (traveling time for P1)	P1B (service time for P1)
Minimum value	2 seconds	13 seconds
Maximum value	10 seconds	55 seconds
Number of data	65 vehicles	65 vehicles
Average value	3.66 seconds	35.85 seconds
Mod value	3 seconds	42 seconds
Standard deviation	1.45 seconds	11.3 seconds
Table 4.5:	Analysis for point inspection P5	– Suspension Test

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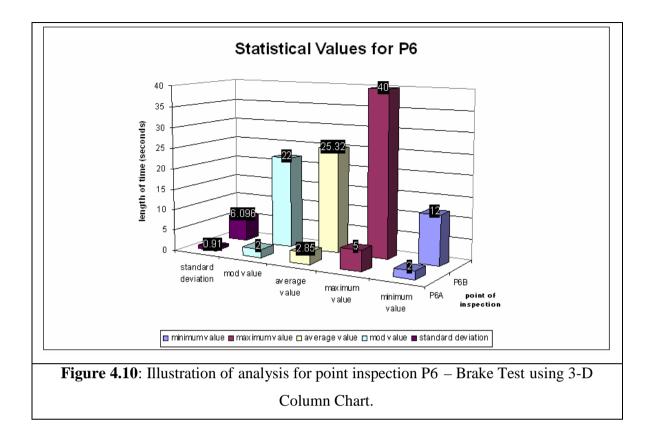
<u>P6 : Brake Test</u>

The analysis for this point of inspection is described in **Table 4.6** below, while the illustration of the figures is shown in **Figure 4.10**.

	P1A (traveling time for P1)	P1B (service time for P1)
Minimum value	2 seconds	12 seconds
Maximum value	5 seconds	40 seconds
Number of data	65 vehicles	65 vehicles
Average value	2.85 seconds	25.32 seconds
Mod value	2 seconds	22 seconds
Standard deviation	0.91 seconds	6.096 seconds
Table 4	4.6: Analysis for point inspection	P6 – Brake Test

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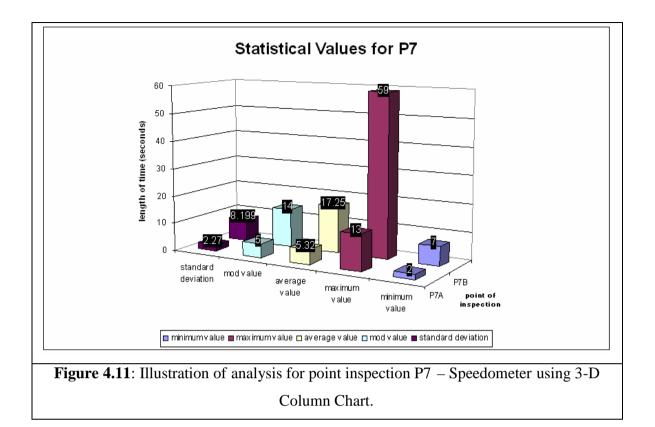
<u>P7 : Speedometer</u>

The analysis for this point of inspection is described in **Table 4.7** below, while the illustration of the figures is shown in **Figure 4.11**.

	P1A (traveling time for P1)	P1B (service time for P1)
Minimum value	2 seconds	7 seconds
Maximum value	13 seconds	59 seconds
Number of data	65 vehicles	65 vehicles
Average value	5.32 seconds	17.25 seconds
Mod value	5 seconds	14 seconds
Standard deviation	2.27 seconds	8.199 seconds
Table 4	7: Analysis for point inspection I	27 – Speedometer

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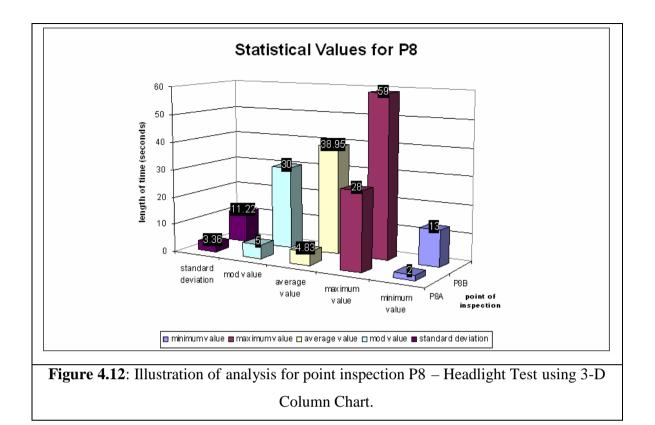


<u>P8 : Headlight Test</u>

The analysis for this point of inspection is described in **Table 4.8** below, while the illustration of the figures is shown in **Figure 4.12**.

	P1A (traveling time for P1)	P1B (service time for P1)
Minimum value	2 seconds	13 seconds
Maximum value	28 seconds	59 seconds
Number of data	65 vehicles	65 vehicles
Average value	4.83 seconds	38.95 seconds
Mod value	5 seconds	30 seconds
Standard deviation	3.36 seconds	11.22 seconds
Table 4.8	Analysis for point inspection P8	8 – Headlight Test

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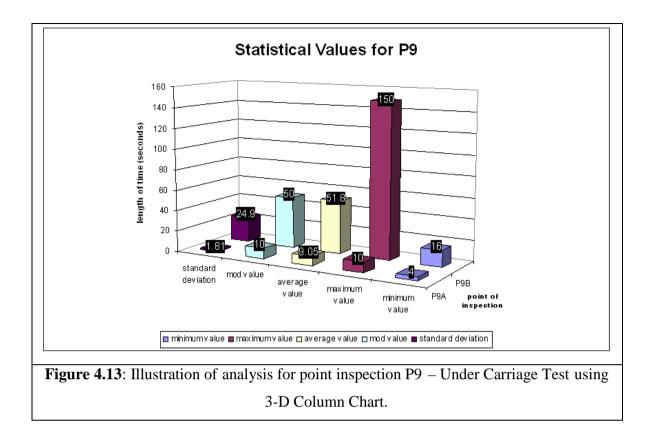


<u>P9 : Undercarriage Check</u>

The analysis for this point of inspection is described in **Table 4.9** below, while the illustration of the figures is shown in **Figure 4.13**.

	P1A (traveling time for P1)	P1B (service time for P1)
Minimum value	4 seconds	16 seconds
Maximum value	10 seconds	150 seconds
Number of data	65 vehicles	65 vehicles
Average value	9.05 seconds	51.8 seconds
Mod value	10 seconds	50 seconds
Standard deviation	1.81 seconds	24.9 seconds
Table 4.9: A	nalysis for point inspection P9 –	Under Carriage Check

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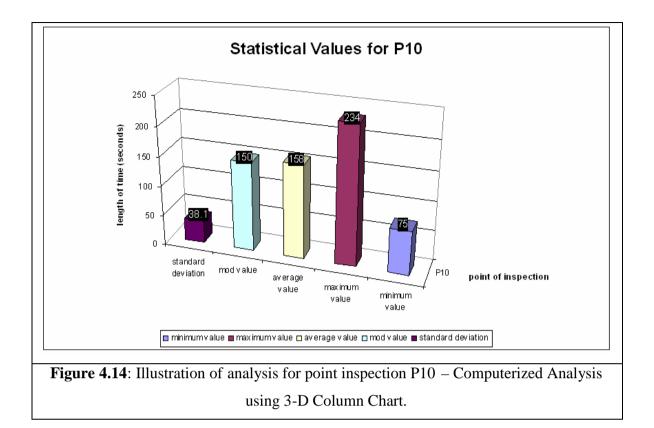


P10 : Computerized Analysis

The analysis for this point of inspection is described in **Table 4.10** below, while the illustration of the figures is shown in **Figure 4.14**.

	P1B (service time for P1)
Minimum value	75 seconds
Maximum value	234 seconds
Number of data	65 vehicles
Average value	158 seconds
Mod value	150 seconds
Standard deviation	38.1 seconds

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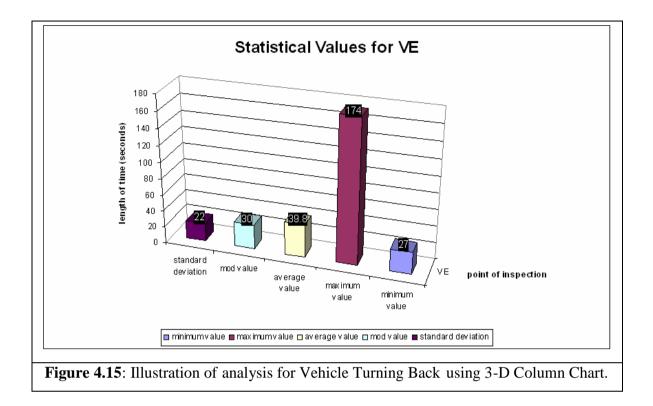


Vehicle Examiner Turning Back

The analysis for situation of vehicle turning back is described in **Table 4.11** below, while the illustration of the figures is shown in **Figure 4.15**.

	P1B (service time for P1)
Minimum value	27 seconds
Maximum value	174 seconds
Number of data	65 vehicles
Average value	39.8 seconds
Mod value	30 seconds
Standard deviation	22 seconds
Table 4.11: Analyse	sis for Vehicle Turning Back

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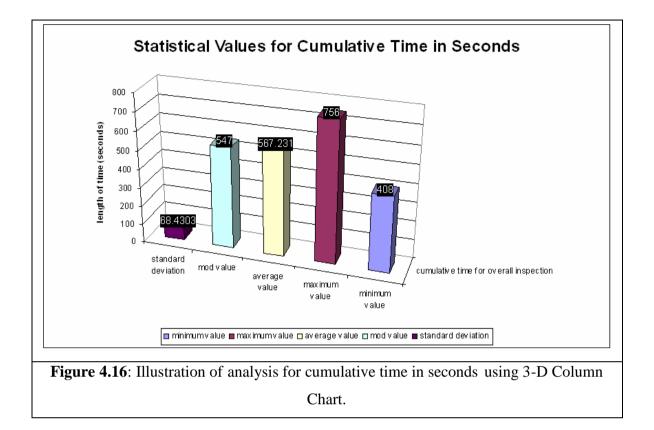


Cumulative Time in Seconds

The analysis for cumulative time in seconds is described in **Table 4.12** below, while the illustration of the figures is shown in **Figure 4.16**.

	P1B (service time for P1)
Minimum value	408 seconds
Maximum value	756 seconds
Number of data	65 vehicles
Average value	567.231 seconds
Mod value	547 seconds
Standard deviation	68.4303 seconds
Table 4.12: Analysis f	for Cumulative Time in Seconds

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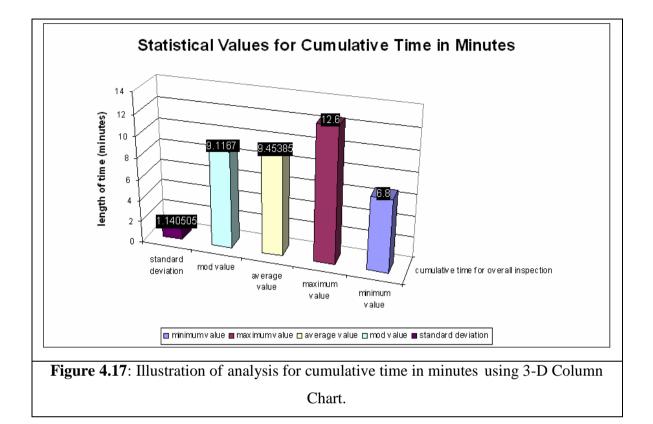


Cumulative Time in Minutes

The analysis for cumulative time in minutes is described in **Table 4.13** below, while the illustration of the figures is shown in **Figure 4.17**.

	P1B (service time for P1)
Minimum value	6.8 seconds
Maximum value	12.6 seconds
Number of data	65 vehicles
Average value	9.45385 seconds
Mod value	9.1167 seconds
Standard deviation	1.140505 seconds
Table 4.13: Analysis	for Cumulative Time in Minutes

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4.4.1 Point Average Summary

The average value for each step of inspection is calculate to get the exact figure on how much time is taken by every inspection for one finance vehicle. The details of the figures are described below, and are shown in **Figure 4.18**.

i.	Average value for traveling time of P1	: 5.91 seconds
ii.	Average value for service time for P1	: 85.14 seconds
iii.	Average value for traveling time of P2	: 4.95 seconds
iv.	Average value for service time for P2	: 2.938 seconds
v.	Average value for traveling time of P3	: 4.71 seconds
vi.	Average value for service time for P3	: 39.52 seconds
vii.	Average value for traveling time of P4	: 3.29 seconds
viii.	Average value for service time for P4	: 28 seconds

ix. Average value for traveling time of P5 :3.66 seconds

x.	Average value for service time for P5	: 35.85 seconds
xi.	Average value for traveling time of P6	: 2.85 seconds
xii.	Average value for service time for P6	:25.32 seconds
xiii.	Average value for traveling time of P7	: 5.32 seconds
xiv.	Average value for service time for P7	: 17.25 seconds
XV.	Average value for traveling time of P8	: 4.83 seconds
xvi.	Average value for service time for P8	: 38.95 seconds
xvii.	Average value for traveling time of P9	: 9.05 seconds
xviii.	Average value for service time for P9	: 51.8 seconds
xix.	Average value for service time for P10	: 158 seconds
XX.	Average value for service time for VE	: 39.8 seconds
xxi.	Average value for service time for cumulation	ive in seconds : 567.231

xxii. Average value for service time for cumulative in minutes : 9.45385

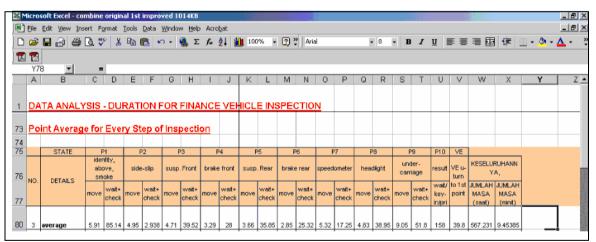


Figure 4.18: Point Average Summary

4.5 Data Simplifying - Approximates Value

Based from the point average summary, the value are rounded and approximate into a range of maximum five seconds plus, as depicted in **Figure 4.19**. This step is done in order to ensure that the process of data analysis being in accurate way and the result will be produced exactly as desired from this research.

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Figure 4.19: A Sample of Data with Approximate Value.

As a result from the above step, the value for each point of inspection detail is listed below and **Figure 4.20** illustrate the result of the analysis.

- i. Approximate value for P1 : 95 seconds
- ii. Approximate value for P2 : 10 seconds
- iii. Approximate value for P3 : 45 seconds
- iv. Approximate value for P4 : 35 seconds
- v. Approximate value for P5 : 40 seconds
- vi. Approximate value for P6 : 30 seconds
- vii. Approximate value for P7 : 25 seconds
- viii. Approximate value for P8 : 45 seconds

- ix. Approximate value for P9 : 65 seconds
- x. Approximate value for P10 : 160 seconds
- xi. Approximate value for VE : 40 seconds
- xii. Approximate value for total inspection time : 590 seconds

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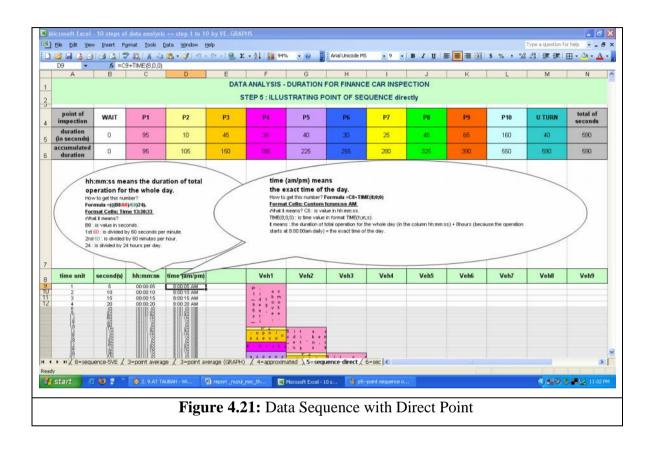
Based form the illustration, it is knew that there is a little difference with the length of cumulative time between the averages for points of inspection which is equal to 568 seconds, while the approximate value for points of inspection is 590 seconds. The total difference between the two values is 22 seconds; this is less than half a minute. So, we assumed the slightly difference in this step as a nil value in this stage of analysis for this study.

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4.6 Data Sequence with Direct Point

The idea for this step of analysis is depicted in the following illustration as shown in **Figure 4.21**.



The two formulas in this step are included in the two columns which are included in these two headers **hh:mm:ss** and **time (am/pm)**. The first formula is noted for considering the duration of total operation for the whole day while the second formula is noted to give the exact time of the day.

Header hh:mm:ss

How do we get this number? There is a toolbar to insert certain formula, and this column is filled with **Formula = (((B8/60)/60)24).**

hh:mm:ss means the duration of total operation for the whole day.

How to get this number?

Formula =(((B8/60)/60)24). Format Cells: Time 13:30:33.

Where;

B8 : is value in seconds.
1st 60 : is divided by 60 seconds per minute.
2nd 60 : is divided by 60 minutes per hour.
24 : is divided by 24 hours per day.

Every cell in this column must be set its format as **Format Cells: Time 13:30:33** as the format itself is provided in Ms Excel 2007.

Header time (am/pm)

The header **time (am/pm)** means the exact time of the day.

How to get this number?

Formula =C8+TIME(8;0;0) and Format Cells: Custom h:mm:ss AM.

Where;

C8 : is value in hh:mm:ss.

TIME(8;0;0) : is time value in format **TIME(h;m;s)**.

The header of the column means the total operation duration for the whole day (in the column hh:mm:ss) added with 8 hours (because the operation starts at 8:00:00am daily) is equal to the exact time of the day. **Figure 4.22** illustrate the content of this step of analysis.

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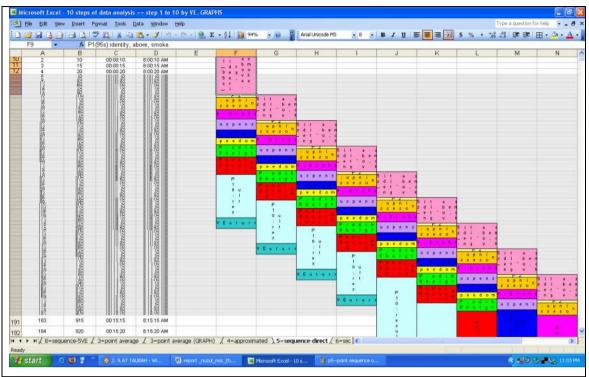


Figure 4.22: Point Sequence Analysis.

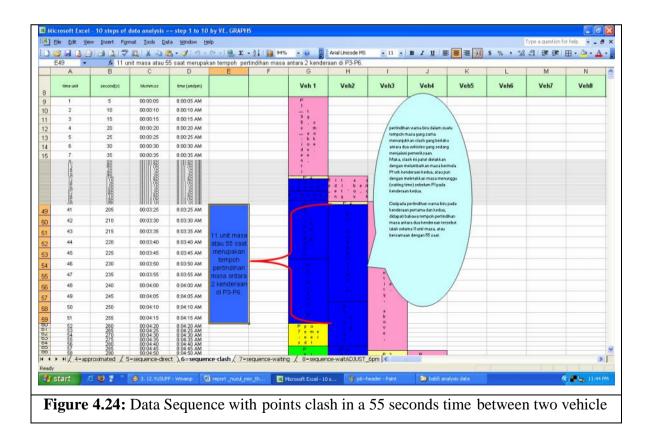
4.7 Data Sequence with Point Clash

In this stage, we gave same color to a few points which being inspected in a same area in the lane of inspection; which is inspected in close proximity between a f ew points of inspection. From the analysis, we colored the third point, P3 to sixth point, P6 as blue cells as they are being inspected in just about the same area in the lane of inspection. The description of this process in depicted in **Figure 4.23**.

Then, we identified the point clash between the two vehicles with same colored points. The blue colored cells show a points clash in a 55 seconds time between two vehicles of Vehicle1 and Vehicle2. See **Figure 4.24**.

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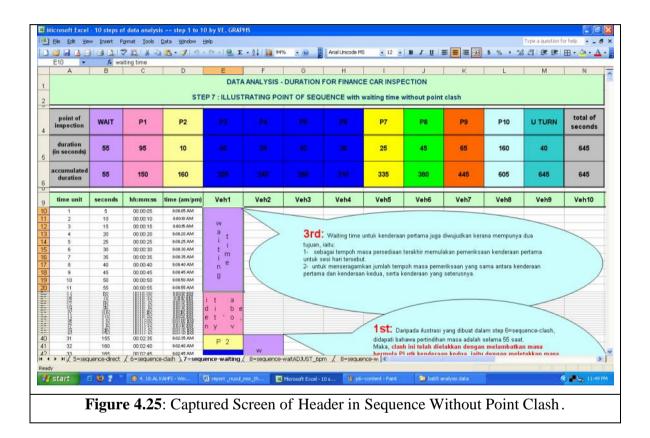
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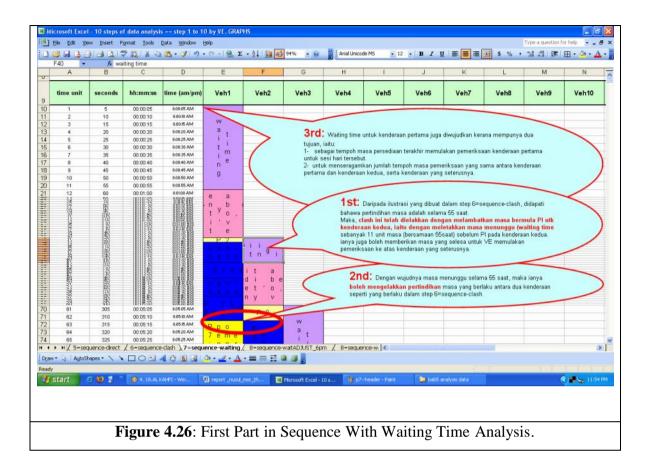
4.8 Data Sequence with Waiting Time and without Point Clash

Knowing that there was some clash between the two series of same colored points, we tried to resolve the situation by adding a period of time called 'waiting time' to make sure that the conflicting points would be separated into different time units. **Figure 4.25** illustrates the heading and the initial content of this analysis. The first purple -colored represents the period of waiting time for each vehicle inspection.



From the previous process, we noticed that there is a point clash in a 55 seconds time. (Refer to **Figure 4.24**). Thus, this clash must be avoided with a waiting time attached to the starting process time (P1) to delay the following vehicle from being conflicting with earlier vehicle. **Figure 4.26** below shows the process of this phase.

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This meaning that a block of 55 seconds waiting time must be allocated to the second vehicles, and also should be given to all vehicles. This waiting time will be able to provide a proper and secure u-turn time for every vehicle examiner to come back from earlier inspected vehicle to the following vehicle that will be inspected. Furthermore, a 55 seconds waiting time (in second callout) will keep away from conflicting the two vehicles as illustrated in **Figure 4.27**.

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As summary, the purposes of waiting time are as follows:

- i. The main point of waiting time is to avoid conflicting between two vehicles in the same time and lane, and of course to make the process of inspection become smoother and comprehensive.
- Waiting time gives a proper time to vehicle examiner come back from the earlier inspected vehicle to the following vehicle that will be inspected.
- iii. Waiting time as a final warm up chance to the first vehicle before starting the inspection.
- iv. Waiting time should be assigned to all vehicles to guarantee that all vehicles in this series obtain a same time period of inspection.

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4.9 Data Sequence with Waiting Time and Adjustment

In this section, we present an adjustment made for the analysis. From the first step illustration in previous process, we also noticed that there is a clash in 9 time units, equally to 45 seconds between the processes of the first and fifth vehicle.

Therefore, this clash must be avoided by putting the fifth vehicle as soon as the first vehicle has completed its last point inspection which means the fifth vehicle will be located after the first VE completed his u-turn from inspecting the first vehicle.

Once the first vehicle has completed its inspection, in next to no time, the fifth vehicle will be carried on with its all points' inspection. This way, there will be no single time unit is unobserved and put in idle situation.

Assigning the fifth vehicle subsequent to the first vehic le, then it will make no problem to allocate the sixth vehicle in no time after the second vehicle without any point clash in any unit time. This second allocation is just keeping up the pattern illustrated in step with waiting time previously.

Why do we need to put the fifth vehicle subsequent to the first vehicle, instead of assigning the sixth or forth vehicle to the first VE? Based from the illustration made in next process as shown in **Figure 4.28**, we have changed the time slot, from vehicles' slot (Veh1, Veh2,...) into vehicle examiners' slot (VE1, VE2, ...). We assume that the step with 4 Vehicle Examiner (VE) is assigning tasks to vehicle examiners in order to complete the inspection process for the day. In slot with waiting time, we noticed that the first vehicle, Veh1 has completed its inspection by the 129th time unit. However, the fifth vehicle is still undergoing its inspection starting from 121st time unit. From the illustration that we have made, we decided to put some time units to wait, bef ore it will be inspected by the first VE as soon as the first vehicle has completed its inspection.

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Figure 4.28: Changing the Header

Therefore, the fifth vehicle is given a period of waiting time as 11 time units as well starting from 121^{st} time unit to 131^{st} time unit. It means the fifth vehicle has to wait the first vehicle completing its inspection entirely before initiating of its inspection can be done. This adjustment is illustrated in **Figure 4.29**.

As a result, we have determined to assigning the fifth vehicle time slot subsequent to the first vehicle's completing time in VE1 (first vehicle examiner) group. Then, we just continue the sixth vehicle being inspected by the VE2 (second vehicle examiner) after completed inspecting second vehicle. This situation is depicted in **Figure 4.30**.

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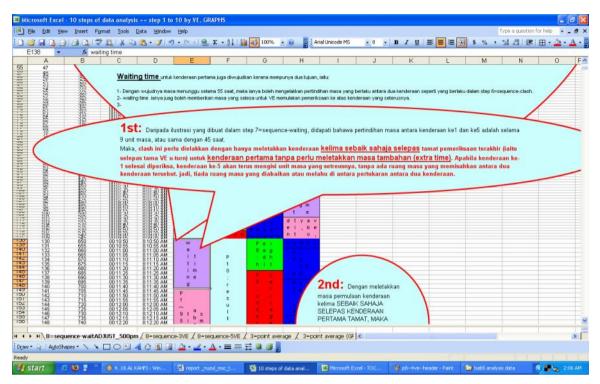
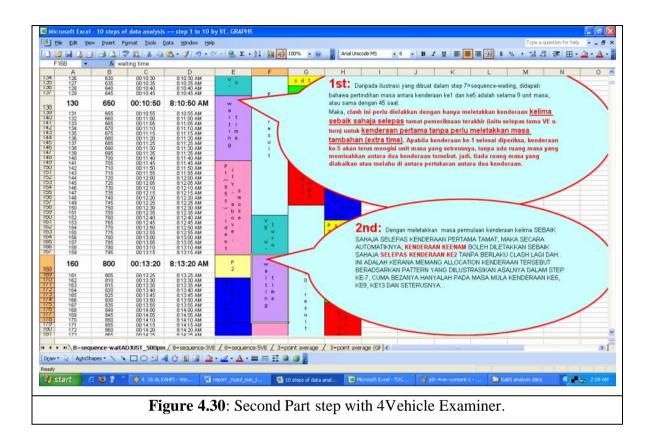


Figure 4.29: First Part step with 4VE.

Why do we set up the fifth vehicle to be inspected subsequent to first vehicle, instead of forth or sixth vehicle? Based on detail justification in the earlier questions, we noticed that the most optimized solution is to allocate the fifth vehicle subsequent to the first vehicle for the first vehicle examiner (and of course not for the fifth vehicle examiner).

In this way, we will be able to control the length of idle time in the first vehicle examiner's slot, which is to achieve our target in order to use all the resources as maximum stage. Specifically, we need to use the lane as well as vehicle examiner at the most optimum point as it can be.

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4.10 Result and Discussion

Based from the analysis, the discussion is divided into 2 cases namely analysis with 3 VE and the other is analysis with 5 VE.

Case 1: Analysis with 3VE

Why not forth vehicle as a restarting point, instead of fifth vehicle? The answer is here. Based on detail explanation presented earlier, we do not assign the forth vehicle, just because to maximize the volume of vehicles can be inspected per day. If we assigned the forth vehicle subsequent to first vehicle, then it will produce less volume of vehicles' inspected. We decide to use the optimum number of four vehicle examiner instead of using only three vehicle examiner at the same time for the reason that we have to maximize the volume of vehicles' inspected for every day. The captured screen for this explanation is shown in **Figure 4.31**.

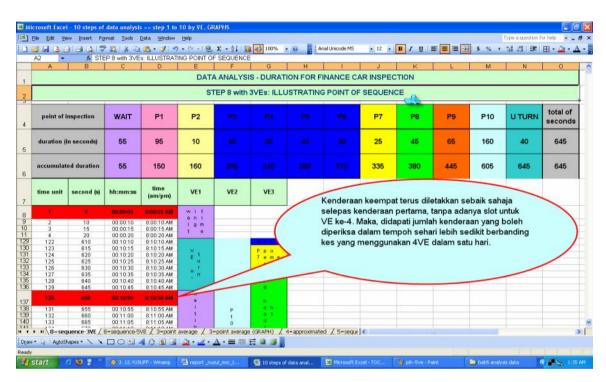


Figure 4.31: The header for the step with waiting time -3VE.

Case 2: Analysis with 5VE

Why not sixth vehicle, instead of fifth vehicle assigned subsequent to first vehicle? In this study, our constraint is to maximize the volume of inspection per day. In the same way as the earlier question, we need to obtain the maximum number of vehicles' inspected produced from the illustration we have made. Therefore, we do not choose to allocate the sixth vehicle subsequent to the first vehic le because of volume constraint. The description of this result is shown in **Figure 4.32**.

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Figure 4.32: The header for the step with waiting time -5VE.

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CHAPTER 5

SYSTEM ARCHITECTURE DESIGN

5.1 Introduction

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Software design is a main part of system engineering. The design elaborated on their behavior process of system with several of diagram. After analyses the user requirement system developer will do the initial design for modeling the new system. It includes the initial plan of design, full design, interface design, database design and so on. One of the important of the design stage is to reduce risk by enabling prediction of system properties. All of this design is covered in this chapter including conceptual design and physical design of ASRS system. The detail process will be shown by use case, activity diagram, sequence diagram and class diagram.

5.2 Current Business Process and Data Model

From observation at vehicle inspection center there are have several problem occur in the current process. Most of the problem in time constrained happed in reservation process. Some of the problem in the as-is system only being able to do very limited operations, such as stated below:

- Customers have to take a long time to wait their vehicle to be inspected.

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- To save the appropriate data of registered vehicles. The vehicles will be registered then they come for initial inspection.
- There is no specific time are allocate to finish for every inspection.
- The system can generate the future appointment with manually by increasing 6 month for every inspection date doing by clerk.
- This system are not identify a holiday. So before create the inspection d ate the clerk will refer to calendar to find the available date.

The figure below showed the overall manual process in vehicle inspection center.

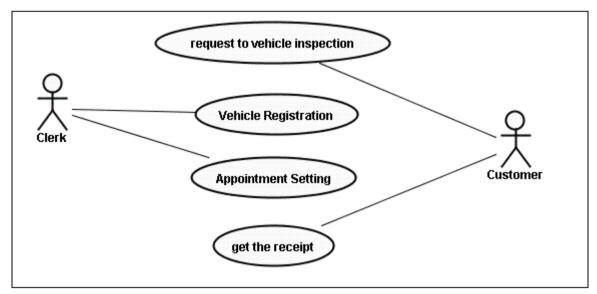


Figure 5.1: Use Case Diagram of Manual Process during Inspection Service

5.3 User Requirements

From the requirement, we know that why the project was proposed, what the problems are who the stakeholders are, what the problem need and what the proposed solution for the new system that we offer to them. A successful project is usually defined as a project that delivers a product which meet user requirement. Similarly, to evaluate the efficiency and success of this process, the result of it need to be comparing against the original requirement.

In order to have better performance on inspection appointment, they required for fully automatic system. As a result, the new system can be more efficient and more effective. In this section describes various functionalities that we propose to be in the new system. The following is a functional requirement requested on new ASRS system.

This user requirement is separated into two sections. First is about Functional Requirement and the second one is Non Functional Requirement. In functional requirement, there are three user involved in this system. There are administ rator, clerk and client. The first section is about user requirement behalf on administrator which handle more on the all function in this system. Clerk can initiate a new and how the registration and reservation process are handled. Finally the last secti on explains how system can provide the services to the customer.

5.3.1 Functional Requirement

Functional requirements capture the intended behavior of the system. This behavior may be expressed as services, tasks or functions the system is required to perform . In product development, it is useful to distinguish between the baseline functionality necessary for any system to compete in that product domain, and *features* that differentiate the system from competitors' products, and from variants in your company's own product line/family. Features may be additional functionality, or differ from the basic functionality along some quality attribute (such as performance or memory utilization). Use cases have quickly become a widespread practice for capturing functional requirements. This is especially true in the object-oriented community where they originated, but their applicability is not limited to object-oriented systems.

Administrator

- 1. To have access for all module for ASRS system i.e. the register
- 2. Have to setting the whole system
- 3. Can update the data in database
- 4. Setting new event in calendar

Clerk and Staff

- 1. To register vehicle data
- 2. To initial a new appointment
- 3. Have searching function to find available date to set the appointment
- 4. Automatic searching function to find the vehicle data by enter the vehicle id
- 5. To view the schedule of appointment
- 6. To print the document and receipt

Client and Customer

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- 1. Be able to view their vehicle information
- 2. Be able to view their appointment date

5.3.2 Non functional requirement

Non-functional requirements are requirements which specify criteria that can be used to judge the operation of a system, rather than specific behaviors. This should be contrasted with functional requirements that specify specific behavior or functions. Non -functional requirements are often called qualities of a system. Other terms for non-functional requirements are constraints, quality attributes, quality goals and quality of service requirements.

1. This system is generated based on online system. So people involve can send or share their information with the centre manager.

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- 2. Provide the efficient scheduling and has an added capability of processing the reservation with automated system.
- 3. Number of lane considered is depending on administrator either one or more. The system will set a default value, but then the admin still can make decision on how many lanes will be use for the day. (Admin can change the value lane for a day, week or month)
- 4. Rule of inspection is that if all three major items passed, the result is pass ed.
- 5. Number of major item is three.
- 6. The three major items named chassis number, engine number and original body.
- 7. Maximum number of vehicle examiner is four.
- 8. Time stagger is one hour (but still not been used in this analysis).
- 9. Time slot is assumed as complete time with waiting time at initial bay.
- Complete time is assumed as processing time with movement period and reserve Time
- 11. Operation hour in this study is defined from 8 a.m. until 5 p.m.
- 12. There is no tea break defined in this study. It is based on the real operation at the time of data been collected.
- 13. There is no lunch hour defined in this study. It is also based on the real operation at the time of data been collected.
- 14. Each vehicle must have a valid time slot of appointment to be inspected. But, if there was any vehicle that came at the wrong date and time, it was also been inspected.
- 15. The system will detect the next appointment by plus six month and it can consider the customer's request and based on any constrained and rule.
- 16. The value of variable can be change for any situation controlled by admin.
- 17. The system will provide searching function for vehicle information and for available slot time.

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18. The searching vehicle information must be flexible. It means the user can define the searching function with any keywords such as by Vehicle Identification Number (VIN), plate number, or class. 19. The function for searching available slot will be generated using formula, and further detail about this formula will be defined in the next section.

5.4 Design Process

Design is a key of the system engineering. In this part, the to -be system will be designed using Unified Modeling Language (UML) as tools to develop a new system. This case, the figure diagram which includes use case diagram, class diagram, sequence diagram and also activity diagram will be illustrated as shown in the next subtopic.

5.4.1 To-Be System Process

Use case Description

Figure 4.2 below represents the whole function of this Appointment Scheduling for Reservation System (ASRS) in vehicle inspection environm ent. This system consist three people invoking the system functionality which are Administrator, clerk and customer.

Clerk	Clerk is a main person who indicates a new appointment. So this					
	system the clerk can do function for insert or update form, setting					
	appointment for insert new appointment and update appointment					
	information, view data vehicle and scheduling appointment and					
	print the schedule.					
Admin	Administrator, people who is responsible for environmental system					
	especially in handling database system and manage the whole					
	system by system setting function. Admin will generate the variable					
	and only admin can change the value of variable then update					
	database.					
Customer	Customer can view the data only including view for data vehicle					
	and the second is view scheduling appointment.					
	Table 5.1. Clarification the Asteria Has Cons					

Table 5.1: Clarification the Actor in Use Case

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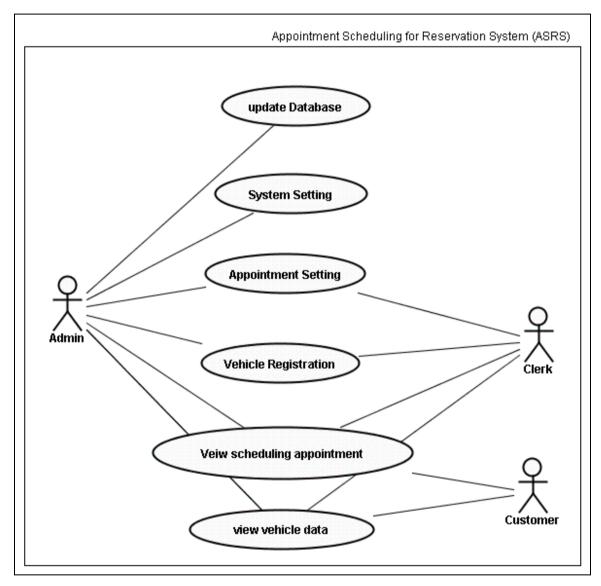


Figure 5.2: Use Case in ASRS System

Use Case Name : Update Database	ID :1	Importance Level: High
Primary Actor : Admin	Use Case	Type : Details, Essential

Stakeholders and Interest :

Admin need to update a new data which related during inspection process. The data including the user account, a new registration for vehicle and the others data

Brief Description :

This use case describes how the admin can update the data in ASRS system. The processes cover on how to edit, add and delete the data.

Trigger: Admin sign up the system first, and achieve the update database function **Type :** External

Relationships:

Association: Include: Extend :

Generalization:

Normal Flow of Events:

- 1. Admin input their username and password for login.
- 2. System verifies the login. If data is valid the update interface will be appear for user in table option.
- 3. Upon successful login, system will display a list of database table have to be updated (S-1)
- 4. System verifies the data entry and validates the data and get the confirmation from user before save to database.
- 5. If user validates, the data will be saving in database permanently.
- 6. The user also can view the latest data and also can print as a document report.

Sub Flows:

S1: Update database process

- 1. Select table data to update
- 2. In this point, admin need to choose either to use search data function or not

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- 3. Then admin can choose the edit, delete or add a new data
- 4. System will be verified the data first before save

Alternate/Exceptional Flows:

Use Case Name : System Setting	ID : 2	Importance Level: High
Primary Actor · Admin	Use Case	Type · Details Essential

Stakeholders and Interest :

Admin – Access System Setting

Brief Description :

This use case, admin need to change the system setting. There are two sub functions which are change value for variable and set up new value for appointment schedu le. That is important to set a new value if any modification had to be in a day.

Trigger: Admin can access the system setting to manage the whole system **Type :** External

Relationships:

Association :

Include :Setting the system, setting schedule, chan ge variable value Extend :

Generalization:

Normal Flow of Events:

- 1. The admin will log in into the system by key in the username and password
- 2. The system will validate whether the username and password valid or not.
- 3. If the login successful, there are have two selection which are change value the variable or term (S-1) and setting schedule function (S-2)

Sub Flows:

S-1: Change value variable

- 1. System views a list of table and a list of variable.
- 2. Admin can change any variable as example to change the term of variable and the initial value for any variable.

S-2:Setting schedule system

- 1. Admin can change the number of lane, number of VE, operation hour, waiting time, working hour
- 2. System will verify the input and save into schedule database

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Alternate/Exceptional Flows:

e Type : Details, Essential
e rype Details, Esselluat
ppointment
or vehicle inspection
ne appointment information
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Use Case Name : view vehicle information	ID :4	Importance Level: High
Primary Actor : Customer	Use Case	Type :Details, Essential
Stakeholders and Interest :		
Customer – view the information only		
Clerk – view the vehicle information after reg	gister the ve	chicle for inspection
Brief Description :		
Customer can view their vehicle information		
Trigger: All user can view this information		
Type : External		
Relationships:		
Association: Clerk		
Include: view vehicle information, pri	int	
Extend :		
Generalization:		
Normal Flow of Events:		
1. The customer log in into the system		
2. System will validate customer login.		
3. Customer need to key in their vehicle		
4. If valid, the appointment information	will be sho	W.
Sub Flows:		
Alternate/Exceptional Flows:		

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Use Case Name : Vehicle Registration	ID :5	Importance Level: High
Primary Actor : Clerk	Use Case	Type : Details, Essential
Stakeholders and Interest :		
Clerk – doing vehicle registration behalf on o	customer	
Brief Description :		
In initial inspection, customer has to inform	for vehicle r	egistration.
Frigger: Clerk will insert the registration for Fype : External	rm.	
Relationships:		
Association : Customer		
Include : vehicle registration, view ve	ehicle data	
Normal Flow of Events:		
1. Login the system by key in username	e and passwo	ord.
2. Registration form will be show		
3. Insert the data in form vehicle. (E.g.		• • • • •
4. System verifies and validates the data		
must to repeat the insert process agai		n.
5. System informs to submit and save the	le	
 The vehicle information will be updated doin The details for this function are given as below The user is logged into the system will login form. System verifies the user login. If data user. User click "Update" button to change User changes the data in form vehicle System verifies and validates the data must to repeat the insert process agai System informs the user to submit ar saving to database. The system ends when the user takes 	ow: The input the a is valid reg ge any inform e. a. If have an n until finish ad save the d	ir username and password via fistration form will be show to th nation. y error in some of data, user h. lata to get confirmation before
the data vehicle.		
the data vehicle.		
the data vehicle. Alternate/Exceptional Flows:		

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Use Case Name : Setting Appointment	ID :6	Importance Level: High
Primary Actor :Clerk	Use Case	Type : Details, Essential
Stakeholders and Interest :		
Clerk - Setting new appointment to vehicle	inspection	
Brief Description :		
The purpose of this function is to set new ap	pointment to	o vehicle inspection and also
update the appointment.		
Trigger: Clerk will initiate a new appointm	ent and also	can update the appointment
Type :External		
Relationships:		
Association : Vehicle data		
Include: Setting appointment, update	e appointmei	nt, view appointment schedule.
Normal Flow of Events:		
1. As usual, clerk must to input their us		
2. System verifies the user login. If dat		erk can choose either to set new
appointment (S-1) or update appoint	tment (S-2)	
Sub Flows:		
S-1 : Set New Appointment		
1. First, user will enter the vehicle num		
2. System checks vehicle record to see		
registered, go to registration process		
3. If the data is available, insert the rec		
System validate the data and check e		•
holiday. If conflict, user must to key		
4. This system also provides the next a		
5. System gets confirmation from user	first before s	save a new data into database.
6. Then fits into the appointment's sch		

S-2: Update Appointment (Request by customer)

- 1. Clerk will enter the vehicle number to determine data of vehicle get appointment information.
- 2. Insert a new date of appointment. System validate the data and check either the data entry are clash or. If conflict, user must to key in another day.

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- 3. System gets confirmation from user first before save a new data into database.
- 4. Then fits into the appointment's schedule.

Alternate/Exceptional Flows:

Activity Diagram

This activity diagram focused on how the customers request the appointment date. When customers come to inspection center for initial inspection, they must to register their vehicle as a record data. The vehicle data will be used in others function in particular to set up the appointment schedule. All the registration and reservation process has been done by clerk behalf on the customer. The inspection will be conducted after all the process finished. At the end of process, the customers get the receipt of the appointment with the inspection number.

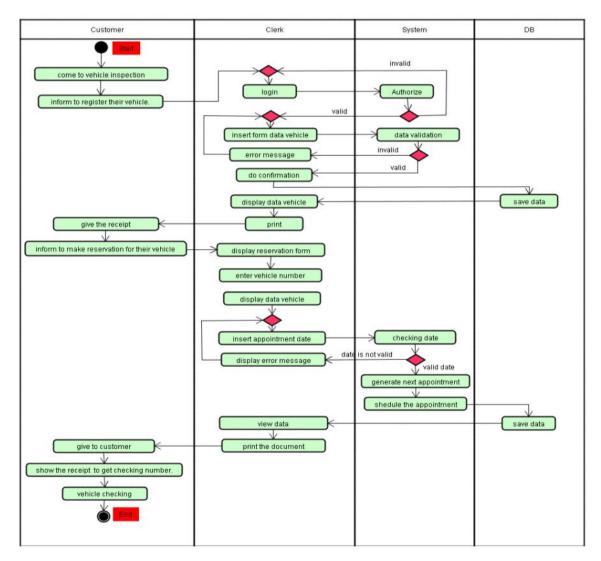


Figure 5.3: Activity Diagram for Vehicle Registration Process

Sequence Diagram

The figures below shows the vehicle registrati on process in propose system. The processes started with login function, followed by insert register form, then system validation to check input error and finally save into database and also can view the whole information. There are three associates class during this process which is Admin class, New Vehicle Class and Vehicle Class.

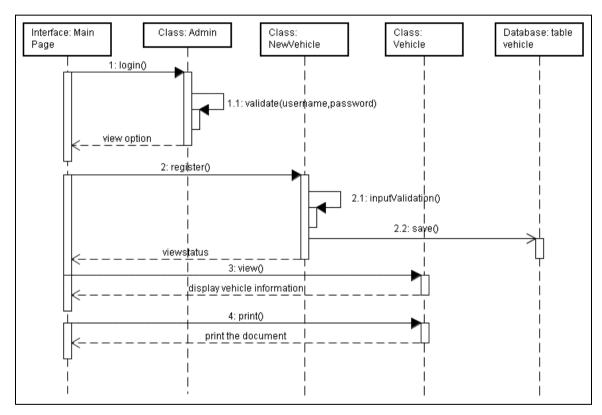


Figure 5.4: Sequence Diagram for Vehicle Registration Process

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This process will be described how clerk setting a new appointment to vehicle inspection behalf on customer. In the end of this function the schedule of the appointment will be generated. This area, the system will generate the propose date automatically based on the available free date from calendar event. The time for inspection will be automatic generated by searching the free slot in reservation schedule.

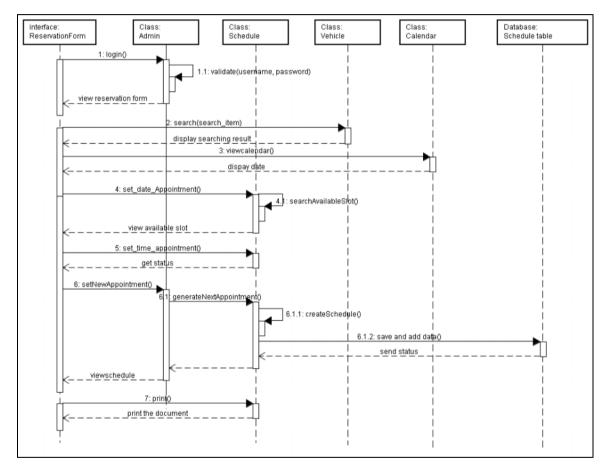


Figure 5.5: Sequence Diagram for Reservation of Appointment Process

DF

The function admin can insert new holiday or event in calendar activity to generate new calendar. The purpose of this function is to find available date before setting any reservation base on the new calendar.

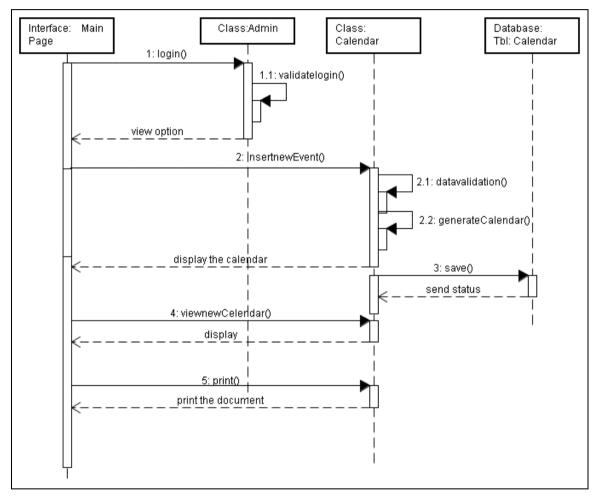


Figure 5.6: Sequence Diagram for Insert New Event Process

SOLID CONVERTER This sequence diagram show the function for update data in database such as change variable, change initial value of data, delete or add a new data. This function admin can select the table of database to update the data.

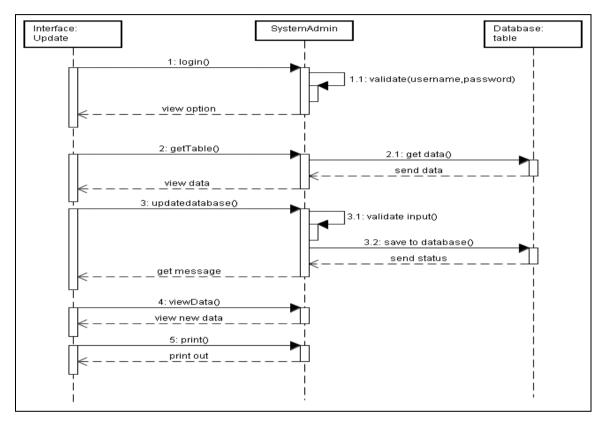


Figure 5.7: Sequence Diagram of Update Table in Database

DF

5.4.2 Conceptual Design

The figure below shows the system architecture for conceptual design of ASRS where all three main user interrelated between them and the ASRS system.

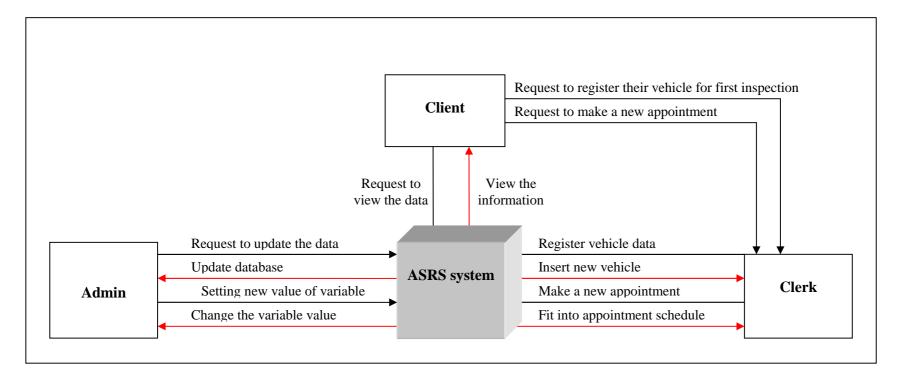


Figure 5.8: Conceptual Design for ASRS System

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5.5 Physical Design

Physical design includes database design, program (structure) chart, interface chart, detailed modules / features and system architecture . Figure 4.9 below shows the whole system architecture for physical design of ASRS system

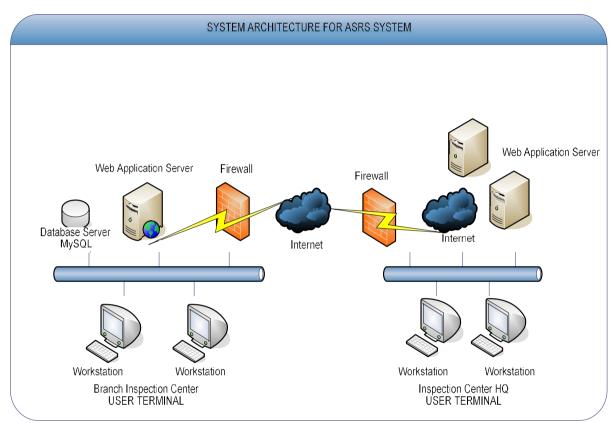


Figure 5.9: Physical Design of ASRS system

5.5.1 Database Design

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Figure 4.10 below shows the class diagram of ARSR system which includes the variable and module for each class. There have eight classes and have some of them had two inherent classes.

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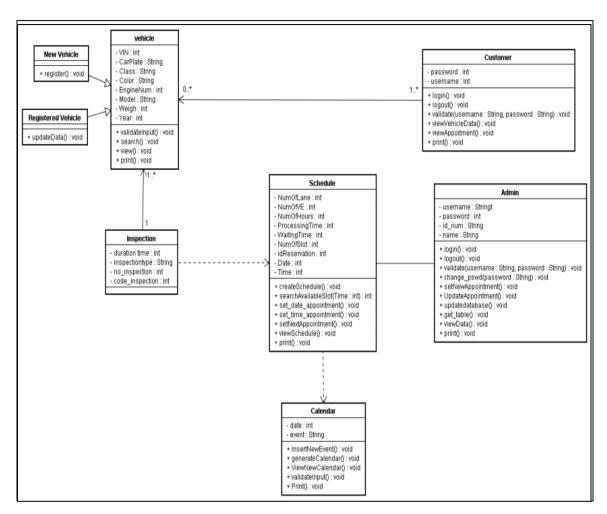


Figure 5.10: Class Diagram of ASRS System

5.5.2 Program Structure

The following Figure 4.11 illustrate the program structure of ASRS system

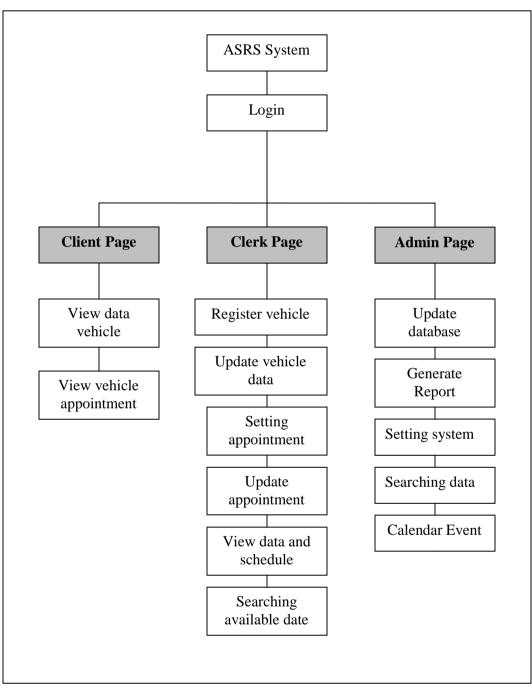


Figure 5.11: ASRS Program Structure

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5.5.3 Interface Chart

This section describes the overall interface chart in ASRS system. There are three main categories which are client page, clerk page and admin page. Table below explain the ASRS features which are assessable to user. The symbol signifies that a specific feature is available to correspondence users.

Features	Admin	Clerk	Client
User registration			
Login page			
Main Page			
About Us			
Inspection Branch			
Administrator main			
System Setting			
Update database			
Setting calendar event			
Update user account			
Searching data			
View vehicle data			
Clerk main page			
View profile			
Vehicle Registration			
Update Appointment			
Setting new appointment			
View appointment information			
Searching available date			

 Table 5.2: ASRS Interface Chart

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CHAPTER 6

HARDWARE AND SOFTWARE REQUIREMENT

6.1 Introduction

This chapter is providing the guide as information resources for hardware and software requirement that required for implement ASRS system in the real environment. This chapter describes the specification of hardware and software requirement to install and run ASRS system. For more detail about how to configure and set up the software, refer to the software installation guides at section 6.4.

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6.2 Hardware requirement

Hardware	Specification	
Windows	Window 98,ME,2000,XP, Vista	
Processor	Intel Pentium 4 and above	
Memory	256 MB and above	
Hard Disk	10 GB and above	
Networking	10/100 Internet	
Printer	Network printer	

The list for hardware requirement is depicted in **Table 5.1**.

Table 6.1: A List of Hardware requirement

6.3 Software requirement

Table 5.2 below shows the list of software requirement that are needed in this project.

Softwa	are Requirement	Purpose
1.	Eclipse Classic 3.3.1.1 Windows:	As a development
	eclipse-SDK-3.3.1.1-win32.zip	platform to write
	http://www.eclipse.org/downloads/	Java programming
2.	Java Development Kit 6 (JDK) update 3:	To install Java
	jdk-6u3-windows-1586-p.exe	Environment
	http://java.sun.com/javase/downloads/index.jsp	
3.	JBossIDE:	
0.	JBossIDE-2.0.0.Beta2-ALL.zip	
	http://labs.jboss.com/jbosside/download/index.html	
	http://https://download/index.itilin_	
4.	JBoss Application Server version 4.2.2.GA:	For Server
	jboss-4.2.2.GA.zip	Application
	http://labs.jboss.com/jbossas/downloads/	11
_		
5.	Database MySQL 5.0.45:	As Server
	mysql-5.0.45-win32.zip	Database
	http://dev.mysql.com/downloads/mysql/5.0.html#win32	
6.	Jboss Control Panel 1.0	For running code
	(given by trainer)	U and
L		

Table 6.2: A List of Software requirement

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6.4 Software Installation

6.4.1 Installation of Java Development Kit 6 (JDK) update 3

• To install JDK, you need to run jdk-6u3-windows-i586-p.exe. You can see the first screen like Figure below. Then you click "**Next**" button.



Figure 6.1: Welcome Screen

• The next screen will be License Agreement. You are advised to read the term of the license carefully before proceeding with the installation. If you are decline the license term, the installation cannot proceed. So click "Accept" button to continue installation.

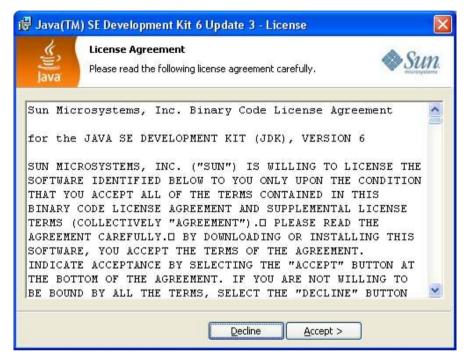


Figure 6.2: License Agreement Screen

At this point you can select any program features for installing Java(TM) SE
Development Kit 6 Update 3. But this project needs you to choose " Development
Tools". At bottom area you allow to change the destination fo lder, if you have specific
reason for doing so. Just click on "Change" button and you can select any folder that you
want. Once you have correct path, click "Next" to proceed.

Select the program features you	u want installed.
Select optional features to install from the list be installation by using the Add/Remove Programs to Development Tools Demos and Samples Source Code Public JRE Java DB	How. You can change your choice of features after utility in the Control Panel Feature Description Java(TM) SE Development Kit 6 Update 3, including private JRE 6 Update 3. This will require 300MB on your hard drive.
ı Install to: C:\Program Files\Java\jdk1.6.0_03\	

Figure 6.3: Custom Setup Screen

SOLID Convert<u>er PDF</u> • Now the program features are being installed. Please wait until the installation finish.

🥵 Java(TM) SE Development Kit 6 Update 3 - Progress			
رچ) Java	Installing The program features you selected are being installed.		
	Please wait while the Install Wizard installs Java(TM) SE Development Kit 6 Update 3. This may take several minutes. Status:		

Figure 6.4: Installing Screen

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• That is Java Setup and need you to choose "Java (TM) SE Runtime Environment". As describe above, you allow changing any destination path for install the Java Program. If you need to use the default setting just continues the next process and click the "Next" button.

🛃 Java Se	tup - Custom	X
(j) Java	Custom Setup Select the program features you want ins	talled.
	M) SE Runtime Environment with support for stures to install from the list below. Java(TM) SE Runtime Environment Additional Languages Default Java for browsers Additional Font and Media Support	European languages will be installed. Select Feature Description The Java(TM) SE Runtime Environment with European languages. This requires 143MB on your hard drive.
Install to: C:\Program	Files\Java\jre1.6.0_03\	
	Cancel	Next >

Figure 6.5: Java Setup Screen

SOLID Conve<u>rter PDF</u> • After all programs are finished, you will see the Wizard completed screen just like Figure below. You can show the readme file to read short information about the installation program. Then click **"Finish"** button to end this process.



Figure 6.6: Complete Screen

• After installation has completed, you required to restart your computer.

6.4.2 Eclipse Installation

Eclipse is an open source community, whose projects are focused on building an extensible development platform, runtimes and application frameworks for building, deploying and managing software across the entire software lifecycle.

The Eclipse open source community has over 60 open source projects. These projects can be conceptually organized into seven different "pillars" or categories:

- 1. Enterprise Development
- 2. Embedded and Device Development
- 3. Rich Client Platform
- 4. Rich Internet Applications
- 5. Application Frameworks
- 6. Application Lifecycle Management (ALM)
- 7. Service Oriented Architecture (SOA)

Refer to the Eclipse website [http://www.eclipse.org] for further information.

• Extract eclipse-SDK-3.3.1.1-win32.zip to C:\. You can select any folder that you want.

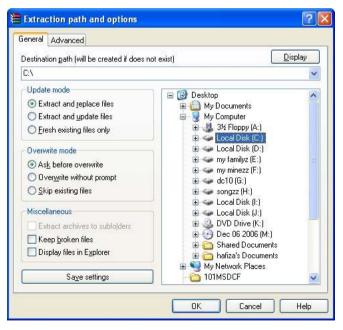


Figure 6.7: Extract eclipse to location folder

• Launch Eclipse application with double click on eclipse.exe

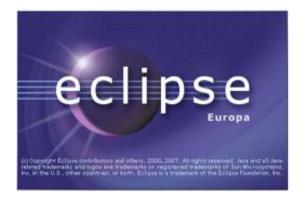


Figure 6.8: Launch Eclipse

• Select path for workspace



Figure 6.9: Select a workspace

• Finish



Figure 6.10: First screen of Eclipse

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6.4.3 JBossIDE Installation

JBoss Eclipse IDE is a set of plug-in designed to be integrated in the Eclipse IDE runtime environment. Some of benefits of JBossIDE are: -

- Extensive and intuitive support for XDoclet.
- The debugging and monitoring of JBoss servers and the controlling of their life cycles.
- An easy way to configure the packaging layout of archives (packed or exploded)
- A simple way to deploy the packaged and/or exploded archive to a JB oss server
- Several J2EE wizards to ease and simplify J2EE development.
- Source code editors for JSP, HTML, and XML

6.4.3.1 Extract JBoss Application Server

• Extract jboss-4.2.2.GA.zip to any location (default is C: \)

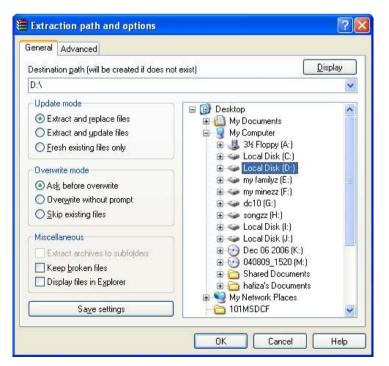


Figure 6.11: Extract Jboss

6.4.3.2 Extract JBoss IDE

• Extract JBossIDE-2.0.0.Beta2-ALL.zip to C:\

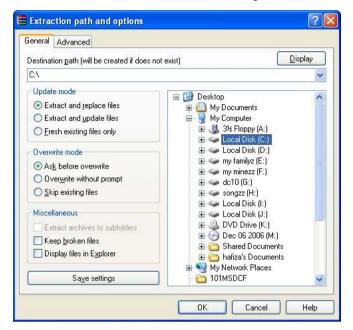


Figure 6.12: Extract Jboss IDE

6.4.3.3 Extract Jboss Control Panel 1.0

• Extract Jboss Control Panel 1.0.zip to any location (default is C: \)

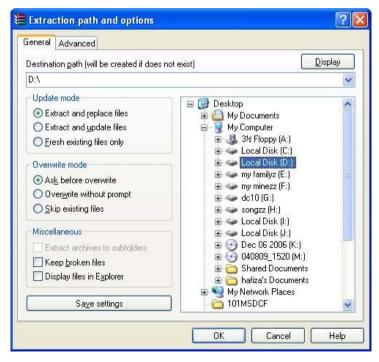


Figure 6.13: Extract Jboss Control Panel 1.0

- Edit config.xml file from Jboss Control Panel 1.0 location:
 - 1 <javaHome>\$path_java_home</javaHome>
 - 2 <locationEAR> \$path_workspace</locationEAR>
 - 3 <jbossHome> \$path_jboss_home</jbossHome>
- Example as **Figure 5.14** below:

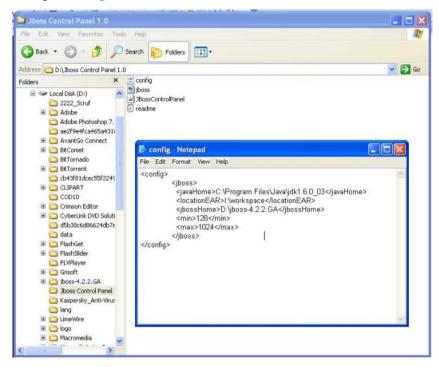


Figure 6.14: Edit config.xml

- To open this control panel, run the JBossControlPanel.jar from JBoss Control Panel 1.0 location.
- The JBoss control panel icon will be displayed in system tray.
- Click on the icon to open the panel.



Figure 6.15: Screen of JBoss Control Panel 1.0

6.4.4 Server Database

In this project, MySQL has been chosen as server database. That is because it gives more advantages. Some of the advantages using MySQL are following: -

- Open source database
- Have fast performance, high reliability and easy to use
- Runs more than 20 platforms including Linux, Windows, OS
- Giving the kind of flexibility

6.4.4.1 MySQL Installation

- Extract mysql-5.0.45-win32.zip to any location
- To run MySQL Server 5.0, double click on "setup.exe" to begin the installation.
- At the welcome screen of setup wizard click "Next" button
- Now you get this window like Figure above. You can choose the type of the installation process. Select the "Complete" option as the setup type then click "Next" button to continue the process



Figure 6.16: Select the setup Type

• Now is ready to install the database into your computer. Click "Install" button to begin installation. Please wait until the installation process finish.

🐻 MySQL Server 5.0 - Setup Wizard	
Ready to Install the Program The wizard is ready to begin installation.	
If you want to review or change any of your installation settings, click Back. Click Cancel exit the wizard. Current Settings:	to
Setup Type: Complete	
Destination Folder: C:\Program Files\MySQL\MySQL Server 5.0\	
< <u>B</u> ack Install Cano	:el

Figure 6.17: Ready to Install MySQL Server

• The MySQL Server Database is now installed. At this point the server must be configured, choose "Configure the MySQLServer now" option and click "Finish" button.

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Figure 6.18: Complete Screen

• Choose a location that will house the database tablespace and click "Next"

MySQL Server Instance Configuration Wizard
MySQL Server Instance Configuration Configure the MySQL Server 5.0 server instance.
Please select the drive for the InnoDB datafile, if you do not want to use the default settings. InnoDB Tablespace Settings Please choose the drive and directory where the InnoDB tablespace should be placed. C: Installation Path Drive Info Volume Name: File System: NTFS IIIIIIIIIIIIIIIIIIIIIIIIIIIIIIIIIII
<pre>< Back Next > Cancel</pre>

Figure 6.19: MySQL Server Instance Configuration

SOLID Conv<u>erter PDF</u> • In this step, you choose enable the root user and choose a password. Check the **"Modify Security Setting"** option, then enter and confirm the new root user password. When finished, click **"Next"**

MySQL Server Instance Configuration	on Wizard	
MySQL Server Instance Configuratio Configure the MySQL Server 5.0 server		\bigcirc
Please set the security options.		
Modify Security Settings		
New root password:	****	Enter the root password.
Confirm:	***	Retype the password.
	🔽 Enable re	oot access from remote machines
Create An Anonymous Account		
This option will create an note that this can lead to		count on this server. Please ystem.
	< Back	Next > Cancel

Figure 6.20: Modify Security Setting

• Now the installation finish. Click "Finish" to close the wizard

MySQL Server Instance Configuration Wizard	×
MySQL Server Instance Configuration Configure the MySQL Server 5.0 server instance.	
Processing configuration	
Prepare configuration Write configuration file (C:\Program Files\MySQL\MySQL Server 5.0\my.ini) Start service	
 ✓ Apply security settings Configuration file created. Windows service MySQL installed. Service started successfully. Security settings applied. Press [Finish] to close the Wizard. 	
< Back Finish Cancel	

Figure 6.21: Finish Screen

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CHAPTER 7

SYSTEM IMPLEMENTATION

7.1 Introduction

In this section, we will discuss briefly the system implementation. Basically there are three levels of user namely High Level Admin, Admin and Clerk.

7.2 High Level Admin

High level admin is referring to those who have the highest authority in the system. It can be the Headquarters office for example the IT Manager or any other post which equivalent to it.

7.2.1 Login

The first user of this system is high level admin. There is only a single position for high level admin in this system. Only one account for high level admin is provided. See **Figure 7.1**. If the user failed to enter the correct login username, the system will prom pt out an error message. These to ensure only the authorized person are allowed to use the system. The source code for this function is shown in **Figure 7.2**.

SCHEDULING.SYSTEM		Welcome
	Username: 0055 Password:	
	Powered by: Nuzulha	
Figure 7.1: Login screen for high level admin.		

String query = "SELECT t1.i_id, t1.i_nama, t1.i_cawangan, t1.i_status, t2.i_nama_cawangan FROM " +
 "puspakom_staf AS t1 JOIN puspakom_cawangan AS t2 ON t2.i_id=t1.i_cawangan " +
 "WHERE t1.i_nostaf=" + username + " AND i_password=" + password + "";
Vector result = db.jdbcOneColumnQuery(query);

Figure 7.2: Source code for login function.

7.2.2 Menu 1: Centre

There a few menus available when we login to the system. The list of menu as depicted in **Figure 7.3** is recorded below:

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- Centre
- Centre Admin
- Public Holidays
- Constraint
- Booking
- Register Inspection
- Start Inspection
- Finish Inspection
- Searching
- Booking Schedule
- Inspection Schedule
- Change Password
- Logout

	SCHEDULING.SYSTEM			
		Welcome		
	Home	III Home		
N	Centre			
2	Centre Admin			
	Public Holidays			
2	Constraint			
8	Booking			
8	Reg. Inspection			
N	Start Inspection			
N	Finish Inspection			
	Searching			
N	Booking Schedule			
	Inspection Schedule			
N	ChangePassword			
	Logout			
	Powered by: Nuzulha			
	Figure 7.3: Menu Screen for super admin.			

A menu for centre management is provided as depicted in **Figure 7.4**. In this section it provides in managing all center's information included address, phone number and fax

number. The high level admin can update and delete the account for corporate office and all centers.

7.2.3 Submenu 1: Add New Centre

Figure 7.5 shows how the system can add new centre. This feat ure can only be manages by high level admin. The values required to add new centre as listed below:

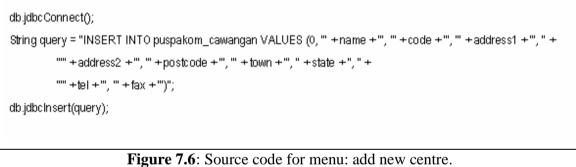
- Centre Name
- Centre Code
- Address
- Postcode
- Town
- State
- Tel. No
- Fax. No

The source code of this feature is shown in Figure 7.6.

	SCHEDULING.SYSTEM Welcome							
	<u>Home</u>	🗄 Cen	tre					
	Centre	[Add New Centre]						
	Centre Admin							
	Public Holidays	No	Centre Name	Tel No	Fax No			
2	Constraint	1	CORPORATE OFFICE	0320527474	0320527427	Þ	\sim	
8	Booking	2	PUSPAKOM JOHOR BAHRU	0712333554	0714325776	Þ	\sim	
2	Reg. Inspection	3	PUSPAKOM KOTA BHARU	097745150	097745153	Þ	\sim × \sim	
	Start Inspection	4	PUSPAKOM WANGSA MAJU	0341433321	0341435388	Þ	\sim	
	Finish Inspection							
	Searching							
	Booking Schedule	oking Schedule						
2	Inspection Schedule							
2	ChangePassword							
8	Logout							
Powered by: Nuzulha								
	Figure 7.4: Centre Menu for super admin.							

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SCHEI	DULING.SYSTEN		elcome				
 Entre Centre Centre Admin Centre Admin Public Holidays Constraint Booking Req. Inspection Start Inspection Start Inspection Searching Booking Schedule Inspection Schedule ChangePassword Logout 	 Add New Centre Centre Name Code Address Postcode Town State Tel. No. Fax No. Submit Reset 	, 					
	Powered by: Nuzulha						
Figure 7.5: Menu for super admin to add new Centre.							



rigure 7.0. Source code for menu. aud new cent

7.2.4 Menu 2: Centre Admin

Figure 7.7 shows the feature on how to manage information on centre admin, or staff of the organization. High level admin can also update and delete the record of admin from other centers. However, the record of high level admin can only be updated.

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The value appeared in this record are listed below:

- Number of staff
- Staff Name

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- Staff Identity Number
- Centre Name
- Mobile phone number

3	Home	⊡C	entre Admin					
9	Centre	[Add New Centre Admin]						
9	Centre Admin							
1	Public Holidays	No				HP No		-
	Constraint	1	NUZULHA KHILWANI HJ IBRAHIM	0055	CORPORATE OFFICE	0197550095	1	
	Booking	2	ABU ALI	1111	PUSPAKOM JOHOR BAHRU	1212121212	P	×
	Reg. Inspection	3	HABIBOLLAH HARON	1133	PUSPAKOM KOTA BHARU	0197777777	D	×
	Start Inspection	4	НАКІМІ	6105	PUSPAKOM KOTA BHARU	0137318499	P	×
	Finish Inspection	5	ROHAIZAN B ABD JALIL	1235	PUSPAKOM WANGSA MAJU		P	×
	Searching	6	NURULHUDA FIRDAUS MOHD AZMI	1122	PUSPAKOM WANGSA MAJU	0193808073	2	×
	Booking Schedule							
	Inspection Schedule							
	ChangePassword							
	Logout							
9								

7.2.5 Submenu1: Add New Centre Admin

High level admin or admin may add new centre admin for their own centre. Please refer to **Figure 7.8**. High level admin can add new record for any new staff in corporate office, while the admin is capable to add new record for any new staff in his centre. The detail information required to add new centre admin is described below:

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- Admin Name
- Staff Number
- No. of Identity Card
- Centre Name
- Mobile phone number

SCHEI	DULING.SYSTEM Welcome				
Home Centre Centre Admin Dublic Holidays Dublic Holidays Constraint Booking Reg. Inspection Start Inspection Searching Booking Schedule Inspection Schedule Inspection Schedule Logout	Add New Centre Admin Admin Name Staff No. NIRC Centre Name -PLEASE SELECT- V H/P No. Submit Reset				
Powered by: Nuzulha					
Figure 7.8: Screen show to add new centre admin.					

The source code to add new centre admin is shown in Figure 7.9.

```
String query = "INSERT INTO puspakom_staf VALUES (0, "' +name +"', "' +staffno +"', " +
""' +password +"', "' +nirc +"', " +centre +", "' +hp +"', 2)";
db.jdbcInsert(query);
```

Figure 7.9: Source code for menu: add new centre admin.

7.3 Managing Public Holidays

The modification of public holidays can be done by both high level admin and normal admin. **Figure 7.10** shows how the process being done. High level admin and admin are capable to add new record for any new public holidays declared for the centre. The detail of information required to add new public holiday is described below:

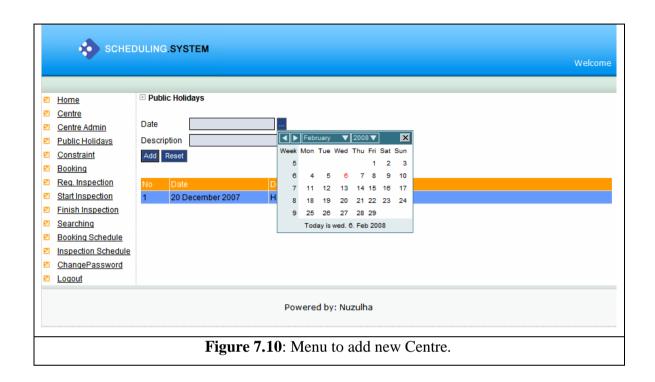
SOLID CONVERTER PDF

• Date

This document was created using

• Description

Then, the list of public holidays will be listed in this screen.



7.4 Constraint

This feature provide on how to set up list of working hours for the centre. Every high level admin or normal admin can make the modification for their own centre. Please refer to **Figure 7.11**.

7.4.1 Submenu 1: Add New Constraint

The submenu that shows in **Figure 7.12** is regards with how to add new working hours for the centre. High level admin and normal admin are capable to add new record for any new working hours declared for the centre.

	SCHED	ULI	NG.S	YSTEM												
															Velc	ome
				-												
2	Home	± 🕻	onstra	Int												
N	<u>Centre</u>	[<u>Ad</u>	d New	Constraint]												
2	Centre Admin															
N	Public Holidays	No		Day			Hours 1				Hours 3		ane I			
N	Constraint					Start		Start				Light	Weight			
2	<u>Booking</u>	1	8	MONDAY	FINANCE / PRIVATE	08:00	13:00	14:00	17:00	null	null	2	0	ALL WEEKS	Þ	\mathbf{X}
N	Reg. Inspection	2	8	MONDAY	COMERCIAL	08:00	13:00	14:00	17:00	null	null	2	2	ALL WEEKS	1	×
2	Start Inspection Finish Inspection	3	8	TUESDAY	FINANCE / PRIVATE	08:00	13:00	14:00	17:00	null	null	2	0	ALL WEEKS	Þ	×
ы	Searching	4	8	TUESDAY	COMERCIAL	08:00	13:00	14:00	17:00	null	null	2	2	ALL WEEKS	Þ	\mathbf{X}
N N	Booking Schedule Inspection Schedule	5	8	WEDNESDAY	FINANCE / PRIVATE	08:00	13:00	14:00	17:00	null	null	2	0	ALL WEEKS	Ď	×
2	ChangePassword	6	8	WEDNESDAY	COMERCIAL	08:00	13:00	14:00	17:00	null	null	2	2	ALL WEEKS	₽	\mathbf{X}
	Logout	7	8	THURSDAY	COMERCIAL	08:00	13:00	14:00	17:00	null	null	2	2	ALL WEEKS	Þ	×
		8	8	THURSDAY	FINANCE / PRIVATE	08:00	13:00	14:00	17:00	null	null	2	0	ALL WEEKS	Þ	×
		9	8	FRIDAY	COMERCIAL	08:00	12:00	15:00	17:00	null	null	2	2	ALL WEEKS	Þ	\mathbf{X}
		10	8	FRIDAY	FINANCE / PRIVATE	08:00	12:00	15:00	17:00	null	null	2	0	ALL WEEKS	Þ	×
		11	8	SATURDAY	FINANCE / PRIVATE	08:00	13:00	null	null	null	null	2	0	ALL WEEKS	Þ	×
	Powered by: Nuzulha															
				Figu	re 7.11: 1	Men	u Lis	t of C	Const	raint	s.					

The detail of information required to add new working hours is described below:

- Duration of processing time
- Type of Day (Monday, or others)
- Type of Week (First Week, or others)
- Working Hour 1 (First block of operation) in time format hh:mm
- Working Hour 2 (Second block of operation) in time format hh:mm
- Working Hour 3 (Third block of operation) in time format hh:mm

SCHEI	DULING.SYSTEM Welcon	ne		
 E Home Centre Centre Admin Public Holidavs Constraint Booking Req.Inspection Start Inspection Finish Inspection Searching Booking Schedule Inspection Schedule ChangePassword Logout 	Add New Constraint Processing Time minutes Day -PLEASE SELECT- Week ALL WEEKS Vehicle Type -PLEASE SELECT- Working Hour 1 to Time Format: HH:mm Working Hour 2 to Time Format: HH:mm Working Hour 3 to Item Time Format: HH:mm Lane Light Submit Reset			
Powered by: Nuzulha				
	Figure 7.12: Menu to add new Centre.			

7.5 Booking Appointment

This is the main menu in the system. This feature provides adding new appointment booking for the centre. Every high level admin or normal admin may add new booking appointment for their own centre. Refer to **Figure 7.13**. High level admin, normal admin and clerk are capable to add new booking for any customer. The detail of information required to add new booking is described below :

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- Centre Name
- Vehicle Category (such as Finance, or others)
- Vehicle Type (Light or Weight)
- Vehicle ID (Plat number)
- Inspection Type (Routine, or others)
- Name of owner
- Phone number of owner
- Address of owner

SCHEI	DULING.SYSTEM	Welcome		
⊠ <u>Home</u>	🗄 Booking			
Centre	Centre Name	PLEASE SELECT		
 <u>Centre Admin</u> <u>Public Holidays</u> 	Vehicle Category			
 <u>Constraint</u> 	Vehicle Type	-PLEASE SELECT-		
Booking	Vehicle ID			
 <u>Req. Inspection</u> <u>Start Inspection</u> 	Trade Plat			
 Start Inspection Finish Inspection 	Inspection Type	PLEASE SELECT		
Searching	Name			
Booking Schedule	Tel No.			
Inspection Schedule ChangePassword	Address			
 ChangePassword Logout 	Submit Reset			
Powered by: Nuzulha				
	Fig	ure 7.13: Menu to Make New Booking.		

The source code of this process is show in Figure 7.14 and Figure 7.15.

7.6 Inspection

This menu provides the feature to add a new inspection for any vehicle that registered at the centre. Every high level admin or normal admin may add new inspection data for their own centre. Refer to **Figure 7.16**. The detail that required adding new inspection is described below:

• Appointment Number

The appointment number can be obtained from the booking slip. Then, the system will register the data recorded in the database and based from it the booking process will be provides.

```
db.jdbcConnect();
String guery = "SELECT * FROM puspakom_constraint WHERE i_cawangan=" + centre +" AND i_veh_type=" +
vehtype;
Vector result = db.jdbcMultipleColumnQuery(query);
if (result.size()!=0){
        while (notbook){
                 query = "SELECT * FROM puspakom_public_holidays WHERE (i_cawangan IS NULL OR
i cawangan="+
                          centre +") AND i_date=" + date[0] +"";
                 result = db.jdbcMultipleColumnQuery(query);
                 if (result.size()!=0) date[1] = "%";
                 query = "SELECT * FROM puspakom_constraint WHERE i_cawangan=" +centre +" AND i_day="
-
                          date[1] +" AND i_week=" +week +" AND i_veh_type=" + vehtype;
                 result = db.jdbcOneColumnQuery(query);
                 if (result.size()==0){
                          query = "SELECT * FROM puspakom_constraint WHERE i_cawangan=" + centre +
                                   " AND i_day=" + date[1] +" AND i_week=6 AND i_veh_type=" + vehtype;
                          result = db.jdbcOneColumnQuery(query);
                 }
      Figure 7.14: First part of the source code to make new appointment booking.
```

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```
for (inti=0;i<lane;i++){
                                 query = "SELECT MAX(i_finish) AS finish FROM puspakom_booking WHERE
                                 i_date="" + date[0] + "' AND i_book_ref=" + (String)result.elementAt(0);
                                 Vector row = db.jdbc One Column Query (query);
                                 if (((String)row.elementAt(0))==null){
                                          row.removeElementAt(0);
                                          row.addElement(s1);
                                 }
                                 row.addElement((String)result.elementAt(0));
                                 finishtime.addElement(row);
                       }
                       ftime = new String[finishtime.size()];
                       fid = new String[finishtime.size()];
                       for (inti=0;i<finishtime.size();i++){
                                 ftime[i] = ((String)((Vector)finishtime.elementAt(i)).elementAt(0)).substring(0, 5);
                                 fid[i] = (String)((Vector)finishtime.elementAt(i)).elementAt(1);
                       }
Figure 7.15: Second part of the source code to make new appointment booking.
```

	SCHEI	DULING.SYSTEM Welcome				
2 2 2 2 2 2 2 2 2 2 2 2 2 2 2 2 2 2 2	Home Centre Centre Admin Public Holidays Constraint Booking Reg. Inspection Start Inspection Start Inspection Searching Booking Schedule Inspection Schedule ChangePassword Logout	Inspection Apmt. Id Submit Reset Without Apmt. Id (Emergency & Special Case Only)				
	Powered by: Nuzulha					
	Figure 7.16: Menu to add new Inspection.					

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7.6.1 Submenu 1: Inspection without Appointment ID and Special Case

For this case of inspection, the staff needs to make new record of inspection to get the details of the vehicle. Please refer to **Figure 7.17**. High level admin, normal admin and clerk are capable to add new booking at the centre. The detail of information required to add new appointment for special case inspection is described below:

- Centre Name
- Vehicle Category (such as Finance, or others)
- Vehicle Type (Light or Weight)
- Vehicle ID (Plat number)
- Inspection Type (Routine, or others)
- Name of owner
- Phone number of owner
- Address of owner

sch	HEDULING.SYSTEM	Welcome			
E Home E Centre E Centre Admin E Public Holidays E Constraint E Booking E Reg. Inspection E Start Inspection E Finish Inspection E Booking Schedule E Booking Schedule E Inspection Schedule E ChangePassword E Logout	ule Submit Reset	-PLEASE SELECT-			
Powered by: Nuzulha					
	Figure 7	17: Submenu to add new Special Inspection			

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7.7 Menu 7: Start and End Inspection

Figure 7.18 record the start time of inspection for every vehicle and **Figure 7.19** record the end time. The detail of information required to record start inspection is described below:

• Queue Number (can be obtained from the start inspection slip).

7.8 Menu 9: Searching

This menu is to search the record of inspection for any vehicle. Refer to **Figure 7.20**. High level admin, normal admin and clerk are capable to search the record of vehicle inspection at the centre. The detail of information required to search informat ion on vehicle inspection is described below:

- Vehicle Number (Plat Number), or
- Trade Plat

7.9 Menu 10: Booking Schedule

This menu is to view the booking schedule for the appropriate date. Refer to **Figure 7.21**. High level admin, normal admin and clerk are capable to view the booking schedule of vehicle inspection at the centre. The detail of information required to view booking schedule is described below:

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- Centre Name
- Date of Inspection

	SCHEE	DULING.SYSTEM Welcome				
2 2 2 2 2 2 2 2 2 2 2 2 2 2 2 2 2 2 2	Home Centre Centre Admin Public Holidays Constraint Bookina Rea, Inspection Start Inspection Finish Inspection Searchina Bookina Schedule	Start Inspection Your Number Submit Reset				
2						
	Figure 7.18: Menu to Record Start Inspection.					

SCHE	DULING.SYSTEM Welcome			
S Home S Centre Centre Admin Public Holidays S Constraint S Booking S Reg. Inspection S Reg. Inspection	Finish Inspection Your Number Submit Reset			
Start Inspection Finish Inspection Searching Booking Schedule Inspection Schedule ChangePassword Logout				
Powered by: Nuzulha				
	Figure 7.19: Menu to Record Finish Inspection.			

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	SCHEL	DULING.SYSTEM Welcom	e		
2	Home	I Searching			
2	Centre				
8	Centre Admin	Vehicle No.			
N	Public Holidays	Trade Plat			
2	Constraint	Search Reset			
2	Booking				
N	Reg. Inspection				
N	Start Inspection				
M	Finish Inspection				
M	Searching				
2	Booking Schedule				
2	Inspection Schedule				
N	ChangePassword				
2	Logout				
	Powered by: Nuzulha				
	Figure 7.20: Menu for Searching.				

SCHEI	DULING.SYSTEM			
	Welco	ome		
Home	Booking Schedule			
Centre				
Centre Admin	Centre NamePLEASE SELECT			
Public Holidays	Date			
Constraint	Search Reset			
Booking				
Req. Inspection				
Start Inspection				
Finish Inspection				
Searching				
Booking Schedule				
Inspection Schedule				
ChangePassword				
Logout				
Powered by: Nuzulha				
	Figure 7.21: Menu to View Booking Schedule.			

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7.10 Menu 11: Inspection Schedule

This menu is to view the inspection schedule for the appropriate date. Refer to **Figure 7.22**. High level admin, normal admin and clerk are capable to view the inspection schedule of vehicle inspection at the centre. The detail of information required to view inspection schedule is described below:

- Centre Name
- Date of Inspection

SCHEI	DULING.SYSTEM	Welcome			
☑ Home	Inspection Schedule				
Centre					
Centre Admin	Centre NamePLEASE SELECT				
Public Holidays	Date				
Constraint	Search Reset				
Booking					
Reg. Inspection					
Start Inspection					
Finish Inspection					
Searching					
Booking Schedule					
Inspection Schedule					
ChangePassword					
Logout					
	Powered by: Nuzulha				
	Figure 7.22: Menu to View Inspection Schedule.				

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7.11 Menu 12: Change Password

This menu is to change the default password for the made for the super admin, admin and clerk. Please refer to **Figure 7.23**. High level admin, normal admin and clerk are capable to change their password for their own accounts. The detail of information required to change password is described below:

- Old Password
- New Password

SCHE	DULING.SYSTEM Welcome			
 E Home Centre Centre Admin Public Holidays Constraint Booking Req.Inspection Start Inspection Searching Booking Schedule Inspection Schedule ChangePassword Logout 	Change Password Password Confirm Password Submit Reset			
Powered by: Nuzulha				
Figure 7.23: Menu to Change Password.				

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CHAPTER 8

CONCLUSION AND DISCUSSION

8.1 Introduction

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This chapter states the conclusion from the discussion of the main finding of this study. It summarizes the overall chapter of this report from Chapter 1 toward Chapter 7. Generally, this chapter has discussed the implication and effect from the project implementation. The following topic is about the constraint and challenging during the project commences and finally the aspiration from the author to the new system proposed.

This report consists of eight main chapters. The eight chapters are Chapter 1 (Introduction), Chapter 2 (Literature Review), Chapter 3 (Methodology), Chapter 4 (Experimental Data and Analysis), Chapter 5 (System Architecture and Design), Chapter 6 (Hardware and Software Requirement), Chapter 7 (System Implementation), and Chapter 8 (Discussion and Conclusion).

Chapter 1 explained about background of the project, statement of project, scope of project, objective of project, the importance of project and chapter summary. In writing a good objective and scope of the project, guidance given from t he project leader is very helpful to ensure the objective have the good quality. A well define scope of project is a key to ensure the success of the project. Following of this is about the case study in vehicle inspection area, with the statement of problem in the current process.

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The literature review which related to the study is then discussed in the Chapter 2. This chapter is divided into 2 main section which the first part is concerning about distributed constraints satisfaction problem and its tech niques that applied in that area. The second part is discussed about software agent and its type in the software environment.

Next, in the Chapter 3 it discussed about the methodology that are been used in the software development. In this chapter, the R UP approach is choose as a methodology for software development. The discussion is not only focusing on the workflow of RUP but also elaborate the detail of its phases such as inception, elaboration, construction and testing phase that are applied in this research. This chapter also discussed about Unified Modeling Language (UML) the software development tools, which includes types of diagrams and notation.

For the purpose of this research, Computerized Vehicle Inspection Centre (PUSPAKOM) has been selected as a real case study. The purpose of selecting PUSPAKOM, is due to the nature of its operation, having many constraints such as capacity of vehicle being inspect for every time slot, processing time, due time, waiting time, release time, availability of working day. The detail explanation of the data collect ion, the experiment and the result from the analysis are discussed in the Chapter 4.

The system architecture and design are thoroughly discussed in Chapter 5. All the process is illustrated in this chapter which included diagrams in UML notation that are use case diagram, sequence diagram, activity diagram and class diagram. The physical and conceptual design of the to-be system is also discussed in this chapter.

In Chapter 6, the hardware and software requirement are discussed. This includes its specification and installation.

For Chapter 7 it is concerning about the system implementation which includes the code design and the implementation of the system. A few snapshot of the system are showed in this chapter.

Finally, chapter 8 provides conclusion and discussion related to the project. The detail the project constraints and challenges are also highlight in this chapter.

8.2 Achievements

Recently, several researchers have applied agent technolog y to such problems especially in manufacturing and business process control. Because of the widely spread lack of theoretically foundation for this kind of collaborative problem solving, the constraint satisfaction/optimization community has recently exte nded the theoretically well-founded constraint satisfaction problem (CSP) model to a distributed constraint satisfaction problems (DCSP). In order to have more effective in appointment scheduling, we introduce a new approach which incorporated constraints that define the constraints network among the agents. As a result, the search algorithm for solving these problems is a distributed algorithm. The algorithm is run by a gents that communicate by sending and receiving messages. In general, messages contain information about assignments of values to variables and refutations of assignments, by agents that have no compatible assignments, to their own variables.

8.3 Constraints

The key constraint is time. Time is very crucial to ensure this project can be achieved. The time given to complete the project is around 1 year and deliverables of this project is conceptual design of the proposed system.

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8.4 Challenges

Completion of the project milestone was one of the challenges. This task was difficult to achieve mainly due to the requirement from the user that needs to be fulfilled.

The challenge was to get real data from the user on the current v ehicle inspection services and business process. This process takes about 2 month that was to gather the original data during the inspection. After the data collection, one of the challenges was to extract and analyze the data and placed it into mathematic al model. This was quite a difficult process because there was a lot of data to be analyzed.

Another challenge was to understand the technique that would be applied in the new proposed system. Then to model the both AS -IS business process and TO-BE business process using UML diagram. By referencing books on UML and attending workshop, the author managed to overcome the challenges to model the data.

8.5 Aspirations

This project aspired to achieved intended objectives agrees earlier in Chapter 1. As a whole, ASRS system has delivered all functionalities stated in the user requirements.

Beside that, other aspiration of this project is to ensure the user is satisfied with the user requirements. By development of this new system, hopefully it can give variou s benefits and make it possible to improve system operation.

Hopefully, this idea of development a new automatic system in vehicle inspection can give various benefits and make it possible to improve system operation, system reliability and to achieve customer and organization satisfactions.

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